

# Automated Planning

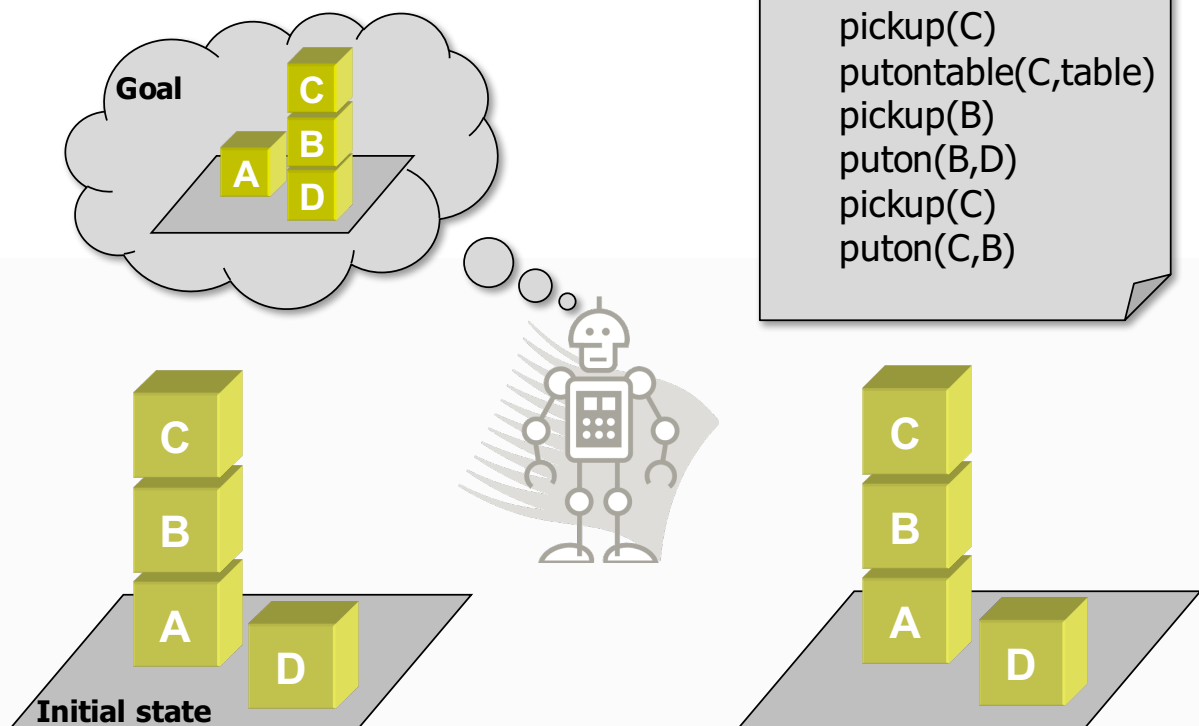
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*What is planning?*

## Blockworld



## **Input:**

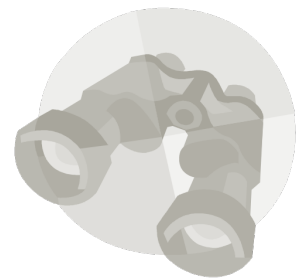
- initial (current) state of the world
- description of actions that can change the world
- desired state of the world

## **Output:**

- a sequence of actions (a plan)

## **Properties:**

- actions in the plan are unknown
- time and resources are not assumed



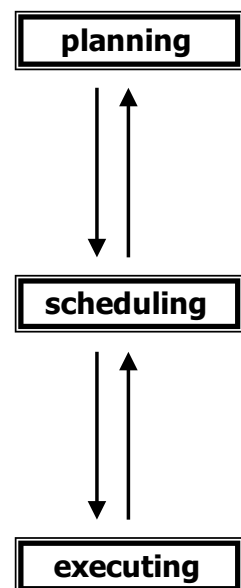
## *Planning and scheduling*

### **Planning**

- deciding which actions are necessary to achieve the goals
- topic of artificial intelligence
- complexity is usually worse than NP-c (in general, undecidable)

### **Scheduling**

- deciding how to process the actions using given restricted resources and time
- topic of operations research
- complexity is typically NP-c





Launch: October 24, 1998

Target: Comet Borrelly

**testing a payload of 12 advanced, high risk technologies**

**— autonomous remote agent**

- planning, execution, and monitoring spacecraft activities based on general commands from operators
- three testing scenarios
  - 12 hours of low autonomy (execution and monitoring)
  - 6 days of high autonomy (operating camera, simulation of faults)
  - 2 days of high autonomy (keep direction)
    - » **beware of backtracking!**
    - » **beware of deadlock in plans!**



*Tutorial outline*

- **Problem Formalisation**
  - models and representations
- **State-space Planning**
  - forward and backward search
- **Plan-space Planning**
  - partial-order planning
- **Control Knowledge in Planning**
  - heuristics
  - control rules



**Planning** deals with **selection and organization of actions** that are changing world states.

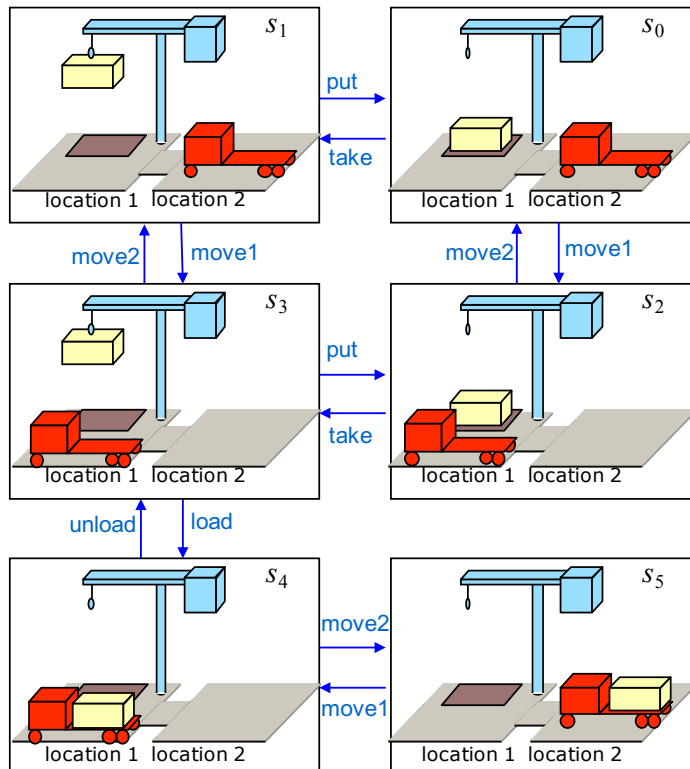
### **System $\Sigma$ modelling states and transitions:**

- **set of states  $S$**  (recursively enumerable)
- **set of actions  $A$**  (recursively enumerable)
  - actions are controlled by the planner!
  - no-op
- **set of events  $E$**  (recursively enumerable)
  - events are out of control of the planner!
  - neutral event  $\varepsilon$
- **transition function  $\gamma: S \times A \times E \rightarrow 2^S$** 
  - actions and events are sometimes applied separately  
 $\gamma: S \times (A \cup E) \rightarrow P(S)$

A planning task is to find which actions are applied to world states to reach some goal from a given initial state.

### **What is a goal?**

- **goal state** or a set of goal states
- **satisfaction of some constraint** over a sequence of visited states
  - for example, some states must be excluded or some states must be visited
- **optimisation of some objective function** over a sequence of visited states (actions)
  - for example, maximal cost or a sum of costs

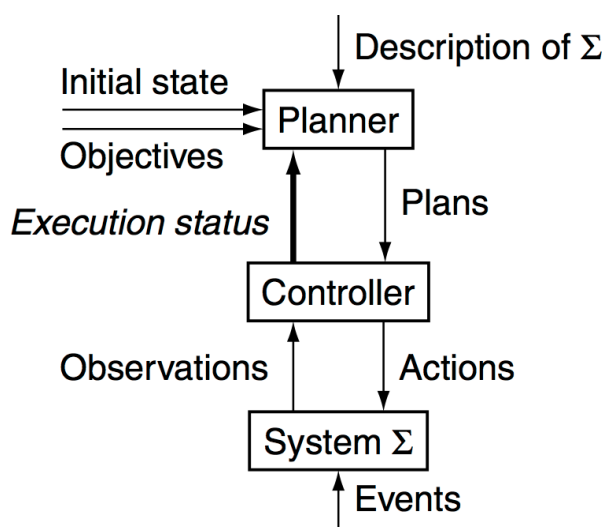


$$\Sigma = (S, A, E, \gamma)$$

- $S = \{s_0, \dots, s_5\}$
- $E = \{\}$  resp.  $\{\varepsilon\}$
- $A = \{\text{move1, move2, put, take, load, unload}\}$
- $\gamma$ : see figure

- init:  $s_0$
- goal:  $s_5$

How does it work?



A **planner** generates plans

A **controller** takes care about plan execution

- for each state it selects an action to execute
- observations (sensor input) are translated to world state

**Dynamic planning involves re-planning when the state is not as expected.**

- the system is **finite**
- the system is **fully observable**
  - We know the current state completely.
- the system is **deterministic**
  - $\forall s \in S \ \forall u \in (A \cup E): |\gamma(s,u)| \leq 1$
- the system is **static**
  - There are no external events.
- the **goals** are **restricted**
  - The aim is to reach one of the goal states.
- the **plans** are **sequential**
  - A plan consists of a (linearly ordered) sequence of actions.
- **time** is **implicit**
  - Actions are instantaneous (no duration is assumed)).
- **planning** is done **offline**
  - State of the world does not change during planning.



We will work with a deterministic, static, finite, and fully observable state-transition system with restricted goals and implicit time  $\Sigma = (S, A, \gamma)$ .

**Planning problem  $P = (\Sigma, s_0, g)$ :**

- $s_0$  is the **initial state**
- $g$  describes the **goal states**

**A solution to the planning problem  $P$**  is a sequence of actions  $\langle a_1, a_2, \dots, a_k \rangle$  with a corresponding sequence of states  $\langle s_0, s_1, \dots, s_k \rangle$  such that  $s_i = \gamma(s_{i-1}, a_i)$  and  $s_k$  satisfies  $g$

👉 **Classical planning (STRIPS planning)** 👉

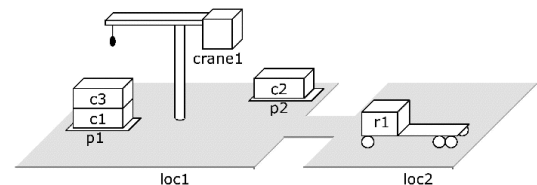
Planning in the restricted model reduces to “path finding” in the graph defined by states and state transitions.

### **Is it really so simple?**

5 locations, 3 piles per location, 100 containers,  
3 robots

↪  **$10^{277}$  states**

This is  $10^{190}$  times more than the largest estimate of the number of particles in the whole universe!



### **How to represent states and actions without enumerating the sets $S$ and $A$ ?**

- recall  $10^{277}$  states with respect to the number of particles in the universe

### **How to efficiently solve planning problems?**

- How to find a path in a graph with  $10^{277}$  nodes?

- **Problem Formalisation**
  - models and representations
- **State-space Planning**
  - Forward and backward search
- **Plan-space Planning**
  - Partial-order planning
- **Control Knowledge in Planning**
  - heuristics
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## Set representation

Each **state** is described using a **set of propositions** that hold at that state.

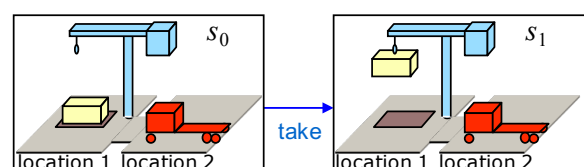
*example: {onground, at2}*

Each **action** is a syntactic expression describing:

- which propositions must hold in a state so the action is applicable to that state  
*example: take: {onground}*
- which propositions are added and deleted from the state to make a new state

*example:*

*take: {onground},  
{holding}<sup>+</sup>*



Let  $L = \{p_1, \dots, p_n\}$  be a finite set of propositional symbols (language).

**A planning domain  $\Sigma$  over  $L$  is a triple  $(S, A, \gamma)$ :**

- $S \subseteq 2^L$ , i.e. **state**  $s$  is a subset of  $L$  describing which propositions hold in that state
  - if  $p \in s$ , then  $p$  holds in  $s$
  - if  $p \notin s$ , then  $p$  does not hold in  $s$
- **action**  $a \in A$  is a triple of subsets of  $L$   
 $a = (\text{precond}(a), \text{effects}^-(a), \text{effects}^+(a))$ 
  - $\text{effects}^-(a) \cap \text{effects}^+(a) = \emptyset$
  - action  $a$  is applicable to state  $s$  iff  $\text{precond}(a) \subseteq s$
- **transition function**  $\gamma$ :
  - $\gamma(s, a) = (s - \text{effects}^-(a)) \cup \text{effects}^+(a)$ , if  $a$  is applicable to  $s$

**Planning problem  $P$  is a triple  $(\Sigma, s_0, g)$ :**

- $\Sigma = (S, A, \gamma)$  is a planning domain over  $L$
- $s_0$  is an initial state,  $s_0 \in S$
- $g \subseteq L$  is a set of goal propositions
  - $S_g = \{s \in S \mid g \subseteq s\}$  is a set of goal states

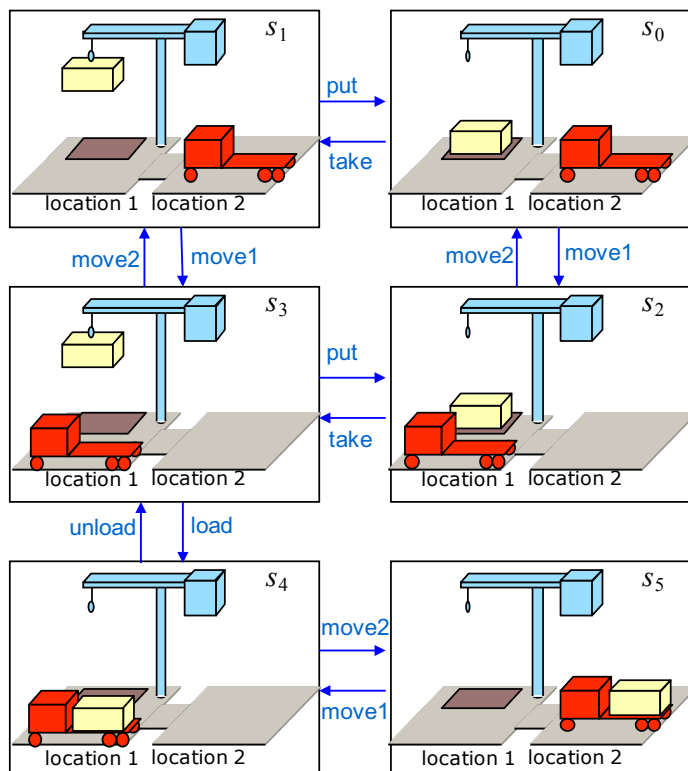
**Plan  $\pi$  is a sequence of actions  $\langle a_1, a_2, \dots, a_k \rangle$**

- **the length of plan  $\pi$**  is  $k = |\pi|$
- **a state obtained by the plan  $\pi$**  (a transitive closure of  $\gamma$ )
  - $\gamma(s, \pi) = s$ , if  $k=0$  (plan  $\pi$  is empty)
  - $\gamma(s, \pi) = \gamma(\gamma(s, a_1), \langle a_2, \dots, a_k \rangle)$ , if  $k>0$  and  $a_1$  is applicable to  $s$
  - $\gamma(s, \pi) = \text{undefined}$ , otherwise

Plan  $\pi$  is a **solution plan** for  $P$  iff  $g \subseteq \gamma(s_0, \pi)$ .

- **redundant plan** contains a subsequence of actions that also solves  $P$
- **minimal plan**: there is no shorter solution plan for  $P$

## Set representation: example



$L = \{\text{onground, onrobot, holding, at1, at2}\}$

$s_0 = \{\text{onground, at2}\}$

$g = \{\text{onrobot}\}$

$\text{load} = (\text{holding, at1}, \text{holding}, \text{onrobot})$

$\langle \text{take, move1, load, move2} \rangle$  is a plan, but not a minimal plan

## Set representation: properties

- Simplicity**

- easy to read

How many states for  $n$  containers?

- Computations**

- the transition function is easy to model/compute using set operations

- if  $\text{precond}(a) \subseteq s$ , then

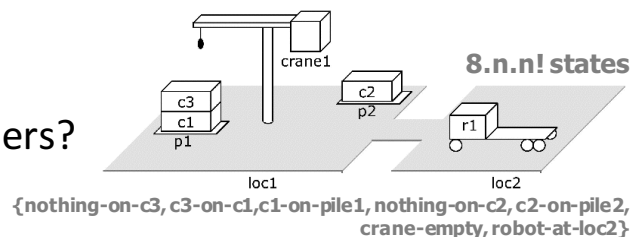
$$\gamma(s, a) = (s - \text{effects}^-(a)) \cup \text{effects}^+(a),$$

- Expressivity**

- some sets of propositions do not describe real states

- $\{\text{holding, onrobot, at2}\}$

- for many domains, the set representation is still too large and not practical



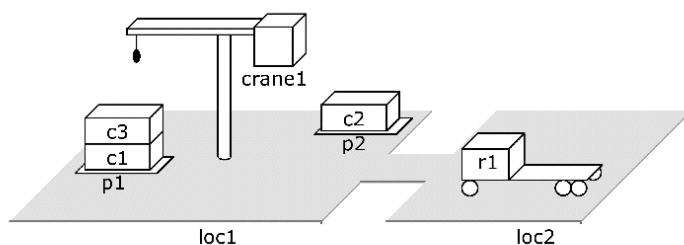
**Classical representation** generalize the set representation by exploiting **first-order logic**.

- **State** is a set of logical atoms that are true in a given state.
- **Action** is an instance of planning operator that changes true value of some atoms.

**More precisely:**

- **L (language)** is a finite set of predicate symbols and constants (there are no function symbols!).
- **Atom** is a predicate symbol with arguments.  
*example:  $on(c3,c1)$*
- We can use **variables** in the operators.  
*example:  $on(x,y)$*

**State is a set of instantiated atoms** (no variables). There is a finite number of states!



{attached(p1,loc1), in(c1,p1), in(c3,p1), top(c3,p1), on(c3,c1), on(c1,pallet), attached(p2,loc1), in(c2,p2), top(c2,p2), on(c2,pallet), belong(crane1,loc1), empty(crane1), adjacent(loc1,loc2), adjacent(loc2,loc1), at(r1,loc2), occupied(loc2), unloaded(r1)}.

- The truth value of some atoms is changing in states:
  - **fluents**
  - *example:  $at(r1,loc2)$*
- The truth value of some state is the same in all states
  - **rigid atoms**
  - *example:  $adjacent(loc1,loc2)$*

We will use a classical **closed world assumption**.

An atom that is not included in the state does not hold at that state!

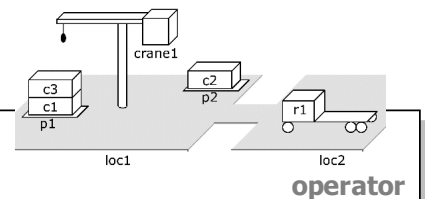
**operator**  $o$  is a triple  $(\text{name}(o), \text{precond}(o), \text{effects}(o))$

- **name( $o$ ):** name of the operator in the form  $n(x_1, \dots, x_k)$ 
  - $n$ : a symbol of the operator (a unique name for each operator)
  - $x_1, \dots, x_k$ : symbols for variables (operator parameters)
    - Must contain all variables appearing in the operator definition!
- **precond( $o$ ):**
  - literals that must hold in the state so the operator is applicable on it
- **effects( $o$ ):**
  - literals that will become true after operator application (only fluents can be there!)

```
take( $k, l, c, d, p$ )  
;; crane  $k$  at location  $l$  takes  $c$  off of  $d$  in pile  $p$   
precond: belong( $k, l$ ), attached( $p, l$ ), empty( $k$ ), top( $c, p$ ), on( $c, d$ )  
effects:  holding( $k, c$ ),  $\neg$ empty( $k$ ),  $\neg$ in( $c, p$ ),  $\neg$ top( $c, p$ ),  $\neg$ on( $c, d$ ), top( $d, p$ )
```

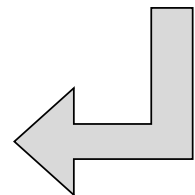
**An action is a fully instantiated operator**

- substitute constants to variables



```
take( $k, l, c, d, p$ )  
;; crane  $k$  at location  $l$  takes  $c$  off of  $d$  in pile  $p$   
precond: belong( $k, l$ ), attached( $p, l$ ), empty( $k$ ), top( $c, p$ ), on( $c, d$ )  
effects:  holding( $k, c$ ),  $\neg$ empty( $k$ ),  $\neg$ in( $c, p$ ),  $\neg$ top( $c, p$ ),  $\neg$ on( $c, d$ ), top( $d, p$ )
```

```
take(crane1,loc1,c3,c1,p1) action  
;; crane crane1 at location loc1 takes c3 off c1 in pile p1  
precond: belong(crane1,loc1), attached(p1,loc1),  
         empty(crane1), top(c3,p1), on(c3,c1)  
effects:  holding(crane1,c3),  $\neg$ empty(crane1),  $\neg$ in(c3,p1),  
          $\neg$ top(c3,p1),  $\neg$ on(c3,c1), top(c1,p1)
```



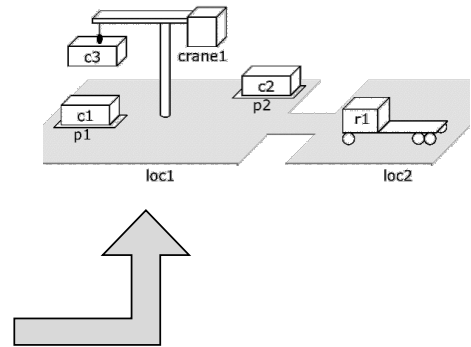
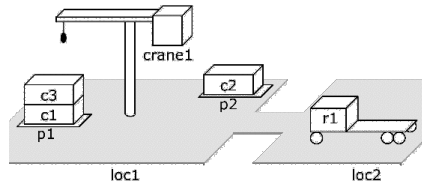
### Notation:

- $S^+ = \{\text{positive atoms in } S\}$
- $S^- = \{\text{atoms, whose negation is in } S\}$

Action **a** is **applicable** to state **s** if and only if  
 $\text{precond}^+(\mathbf{a}) \subseteq \mathbf{s} \wedge \text{precond}^-(\mathbf{a}) \cap \mathbf{s} = \emptyset$

The result of application of action **a** to **s** is  
 $\gamma(\mathbf{s}, \mathbf{a}) = (\mathbf{s} - \text{effects}^-(\mathbf{a})) \cup \text{effects}^+(\mathbf{a})$

```
take(crane1, loc1, c3, c1, p1)
;; crane crane1 at location loc1 takes c3 off c1 in pile p1
precond: belong(crane1, loc1), attached(p1, loc1),
         empty(crane1), top(c3, p1), on(c3, c1)
effects:  holding(crane1, c3), ¬empty(crane1), ¬in(c3, p1),
         ¬top(c3, p1), ¬on(c3, c1), top(c1, p1)
```



Let  $L$  be a language and  $O$  be a set of operators.

**Planning domain  $\Sigma$**  over language  $L$  with operators  $O$  is a triple  $(S, A, \gamma)$ :

- **states**  $S \subseteq 2^{\{\text{all instantiated atoms from } L\}}$
- **actions**  $A = \{\text{all instantiated operators from } O \text{ over } L\}$ 
  - action **a** is **applicable** to state **s** if  
 $\text{precond}^+(\mathbf{a}) \subseteq \mathbf{s} \wedge \text{precond}^-(\mathbf{a}) \cap \mathbf{s} = \emptyset$
- **transition function  $\gamma$** :
  - $\gamma(\mathbf{s}, \mathbf{a}) = (\mathbf{s} - \text{effects}^-(\mathbf{a})) \cup \text{effects}^+(\mathbf{a})$ , if **a** is applicable on **s**
  - $S$  is closed with respect to  $\gamma$  (if  $\mathbf{s} \in S$ , then for every action **a** applicable to **s** it holds  $\gamma(\mathbf{s}, \mathbf{a}) \in S$ )

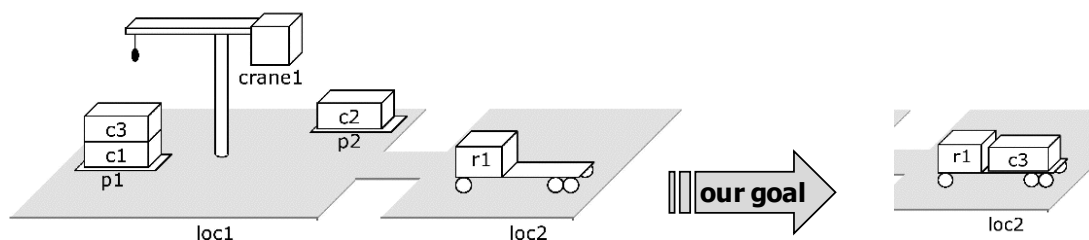
**Planning problem P** is a triple  $(\Sigma, s_0, g)$ :

- $\Sigma = (S, A, \gamma)$  is a planning domain
- $s_0$  is an initial state,  $s_0 \in S$
- $g$  is a set of instantiated literals
  - state  $s$  satisfies the goal condition  $g$  if and only if  $g^+ \subseteq s \wedge g^- \cap s = \emptyset$
  - $S_g = \{s \in S \mid s \text{ satisfies } g\}$  – a set of goal states

Usually the planning problem is given by a triple  $(O, s_0, g)$ .

- $O$  defines the operators and predicates used
- $s_0$  provides the particular constants (objects)

Classical representation: an example plan

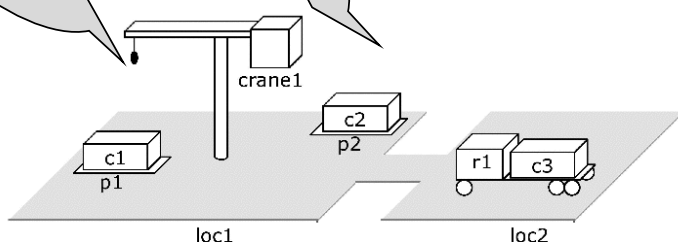


$s_1 = \{ \text{attached}(p1, loc1), \text{in}(c1, p1), \text{in}(c3, p1), \text{top}(c3, p1), \text{on}(c3, c1), \text{on}(c1, pallet), \text{attached}(p2, loc1), \text{in}(c2, p2), \text{top}(c2, p2), \text{on}(c2, pallet), \text{belong}(\text{crane1}, loc1), \text{empty}(\text{crane1}), \text{adjacent}(loc1, loc2), \text{adjacent}(loc2, loc1), \text{at}(r1, loc2), \text{occupied}(loc2), \text{unloaded}(r1) \}.$

$g = \{ \text{loaded}(r1, c3), \text{at}(r1, loc2) \}$

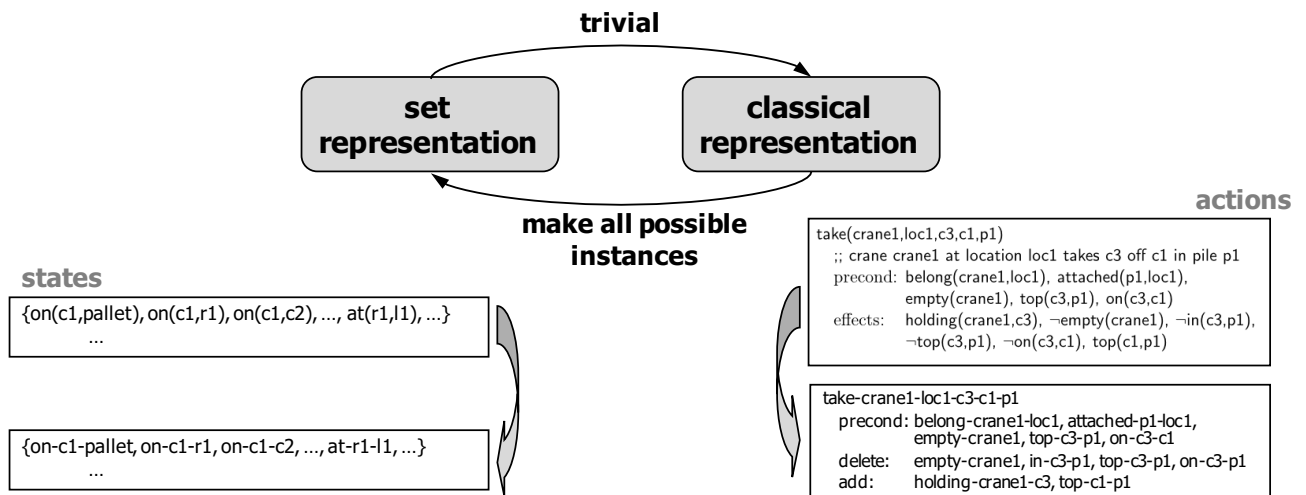
take(crane1, loc1, c3, c1, p1),  
move(r1, loc2, loc1),  
load(crane1, loc1, c3, r1),  
move(r1, loc1, loc2)

move(r1, loc2, loc1),  
take(crane1, loc1, c3, c1, p1),  
load(crane1, loc1, c3, r1),  
move(r1, loc1, loc2)



**Expressive power** of both representations is **identical**.

However, the translation from the classical representation to a set representation brings **exponential increase of size**.

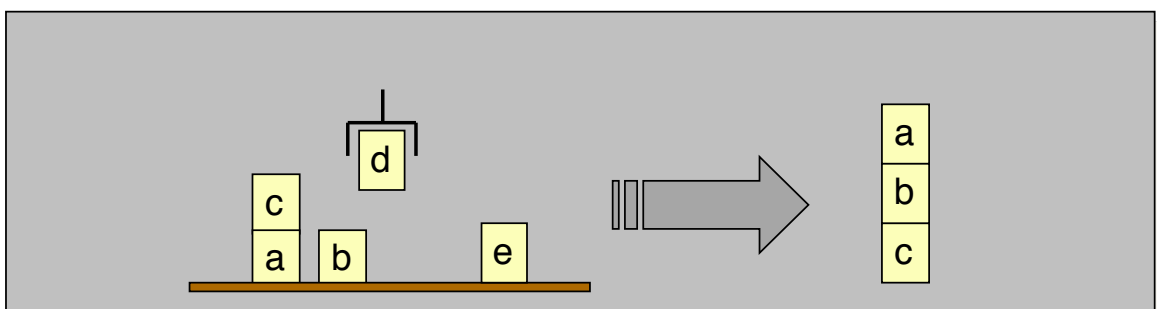


## Blockworld: an example problem

### The blocks world

- infinitely large table with a finite set of blocks
- the exact location of block on the table is irrelevant
- a block can be on the table or on another (single) block
- the planning domain deals with moving blocks by a computer hand that can hold at most one block

*situation example*



## Blockworld: classical representation

### Constants

- blocks: a,b,c,d,e

### Predicates:

- **ontable(x)**  
block x is on a table
- **on(x,y)**  
block x is on y
- **clear(x)**  
block x is free to move
- **holding(x)**  
the hand holds block x
- **handempty**  
the hand is empty

### Actions

#### unstack(x,y)

Precond:  $\text{on}(x,y), \text{clear}(x), \text{handempty}$   
Effects:  $\neg \text{on}(x,y), \neg \text{clear}(x), \text{clear}(y),$   
 $\neg \text{handempty}, \text{holding}(x),$

#### stack(x,y)

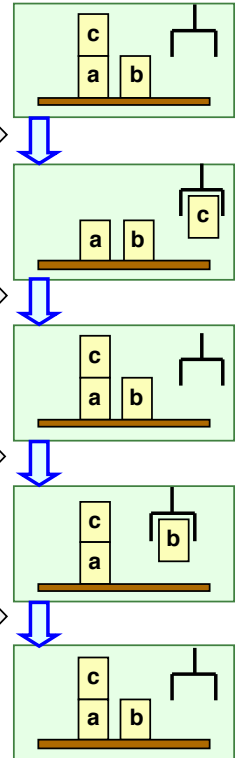
Precond:  $\text{holding}(x), \text{clear}(y)$   
Effects:  $\neg \text{holding}(x), \neg \text{clear}(y),$   
 $\text{on}(x,y), \text{clear}(x), \text{handempty}$

#### pickup(x)

Precond:  $\text{ontable}(x), \text{clear}(x), \text{handempty}$   
Effects:  $\neg \text{ontable}(x), \neg \text{clear}(x),$   
 $\neg \text{handempty}, \text{holding}(x)$

#### putdown(x)

Precond:  $\text{holding}(x)$   
Effects:  $\neg \text{holding}(x), \text{ontable}(x),$   
 $\text{clear}(x), \text{handempty}$



## Blockworld: set representation

### Propositions:

36 propositions for 5 blocks

- **ontable-a**  
block **a** is on table (5x)
- **on-c-a**  
block **c** is on block **a** (20x)
- **clear-c**  
block **c** is free to move (5x)
- **holding-d**  
the hand holds block **d** (5x)
- **handempty**  
the hand is empty (1x)

### Actions

50 actions for 5 blocks

#### unstack-c-a

Pre:  $\text{on-c-a}, \text{clear-c}, \text{handempty}$   
Del:  $\text{on-c-a}, \text{clear-c}, \text{handempty}$   
Add:  $\text{holding-c}, \text{clear-a}$

#### stack-c-a

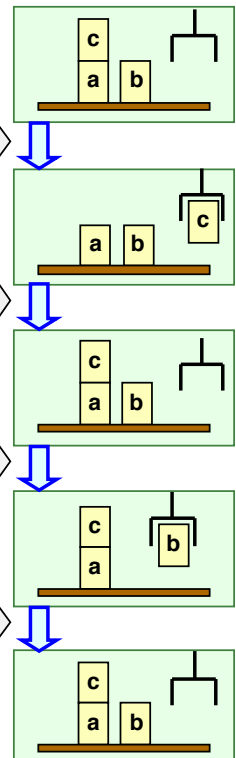
Pre:  $\text{holding-c}, \text{clear-a}$   
Del:  $\text{holding-c}, \text{clear-a}$   
Add:  $\text{on-c-a}, \text{clear-c}, \text{handempty}$

#### pickup-b

Pre:  $\text{ontable-b}, \text{clear-b}, \text{handempty}$   
Del:  $\text{ontable-b}, \text{clear-b}, \text{handempty}$   
Add:  $\text{holding-b}$

#### putdown-b

Pre:  $\text{holding-b}$   
Del:  $\text{holding-b}$   
Add:  $\text{ontable-b}, \text{clear-b}, \text{handempty}$



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## State-space planning

**The search space corresponds to the state space of the planning problem.**

- search nodes correspond to world states
- arcs correspond to state transitions by means of actions
- the task is to find a path from the initial state to some goal state

### Basic approaches

- forward search
- backward search
  - lifting
  - STRIPS
- problem dependent (blocks world)

*Note: all algorithms will be presented for the classical representation*

**Start in the initial state and go towards some goal state.**

We need to know:

- whether a given state is a **goal state**
- how to find a set of **applicable actions** for a given state
- how to define a state after **applying a given action**

### Forward planning: algorithm

Forward-search( $O, s_0, g$ )

$s \leftarrow s_0$

$\pi \leftarrow$  the empty plan

loop

if  $s$  satisfies  $g$  then return  $\pi$

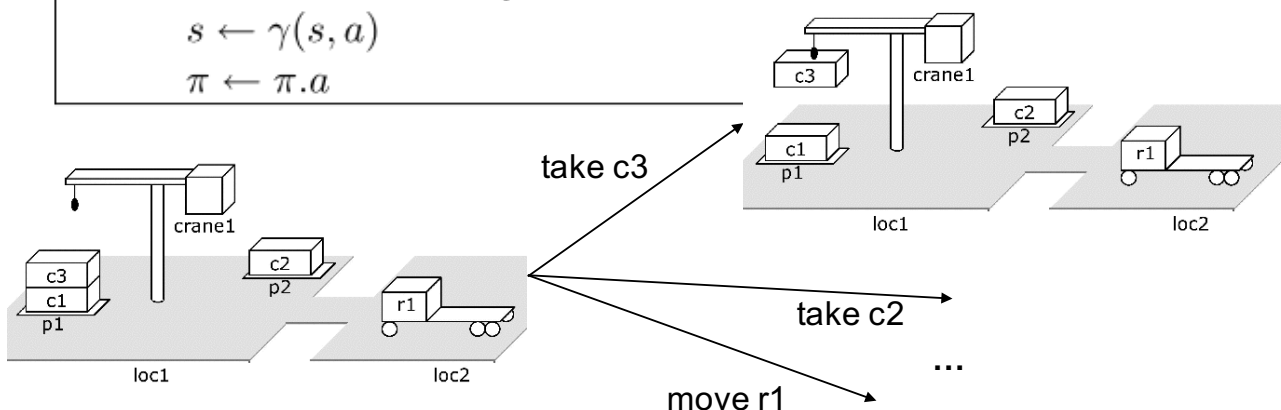
$E \leftarrow \{a \mid a \text{ is a ground instance of an operator in } O, \text{ and } \text{precond}(a) \text{ is true in } s\}$

if  $E = \emptyset$  then return failure

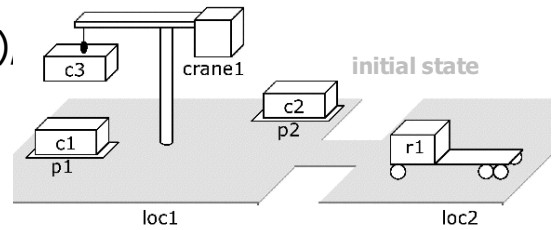
nondeterministically choose an action  $a \in E$

$s \leftarrow \gamma(s, a)$

$\pi \leftarrow \pi.a$



{belong(crane1,loc1), adjacent(loc2,loc1),  
holding(crane1,c3), unloaded(r1),  
at(r1,loc2),  $\neg$ occupied(loc1),  
occupied(loc2),...}



**move(r1,loc2,loc1)**

move( $r, l, m$ )  
;; robot  $r$  moves from location  $l$  to location  $m$   
precond: adjacent( $l, m$ ), at( $r, l$ ),  $\neg$ occupied( $m$ )  
effects: at( $r, m$ ), occupied( $m$ ),  $\neg$ occupied( $l$ ),  $\neg$ at( $r, l$ )

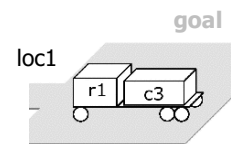
{belong(crane1,loc1),  
adjacent(loc2,loc1), holding(crane1,c3), unloaded(r1),  
at(r1,loc1), occupied(loc1), ...}

**load(crane1,loc1,c3,r1)**

load( $k, l, c, r$ )  
;; crane  $k$  at location  $l$  loads container  $c$  onto robot  $r$   
precond: belong( $k, l$ ), holding( $k, c$ ), at( $r, l$ ), unloaded( $r$ )  
effects: empty( $k$ ),  $\neg$ holding( $k, c$ ), loaded( $r, c$ ),  $\neg$ unloaded( $r$ )

{belong(crane1,loc1), adjacent(loc2,loc1),  
empty(crane1), loaded(r1,c3),  
at(r1,loc1), occupied(loc1), ...}

**Goal = {at(r1,loc1),loaded(r1,c3)}**



## Forward planning algorithm is sound.

- If some plan is found then it is a solution plan..
- It is easy to verify by using  $s = \gamma(s_0, \pi)$ .

## Forward planning algorithm is complete.

- If there is any solution plan then at least one search branch corresponds to this plan.
- induction by the plan length
- at each step, the algorithm can select the correct action from the solution plan (if correct actions were selected in the previous steps)

We need to implement the presented algorithm in a deterministic way:

– **breadth-first search**

- sound, complete, but memory consuming

– **depth-first search**

- sound, completeness can be guaranteed by loop checks (no state repeats at the same branch)

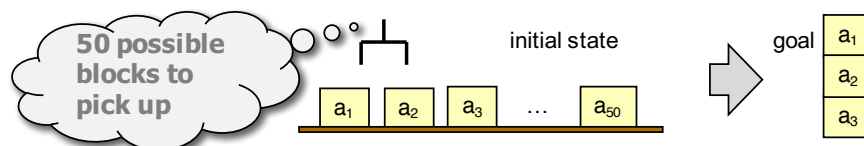
– **A\***

- if we have some admissible heuristic
- the most widely used approach

*Branching*

What is the major problem of forward planning?

**Large branching factor** – the number of options



- This is a problem for deterministic algorithm that needs to explore the possible options.

**Possible approaches:**

- **heuristic** recommends an action to apply
- **pruning** of the search space

- For example, if plans  $\pi_1$  and  $\pi_2$  reached the same state then we know that plans  $\pi_1 \pi_3$  and  $\pi_2 \pi_3$  will also reach the same state. Hence the longer of the plans  $\pi_1$  and  $\pi_2$  does not need to be expanded. We need to remember the visited states ☹.

**Start with a goal (not a goal state as there might be more goal states) and through sub-goals try to reach the initial state.**

We need to know:

- whether the state **satisfies the current goal**
- how to find **relevant actions** for any goal
- how to define the **previous goal** such that the action converts it to a current goal

### Backward planning: relevant actions

**Action  $a$  is relevant for a goal  $g$  if and only if:**

- action  $a$  contributes to goal  $g$ :  $g \cap \text{effects}(a) \neq \emptyset$
- effects of action  $a$  are not conflicting goal  $g$ :
  - $g^- \cap \text{effects}^+(a) = \emptyset$
  - $g^+ \cap \text{effects}^-(a) = \emptyset$

A **regression set** of the goal  $g$  for (relevant) action  $a$  is

$$\gamma^{-1}(g,a) = (g - \text{effects}(a)) \cup \text{precond}(a)$$

**Example:**

goal:  $\{\text{on}(a,b), \text{on}(b,c)\}$

action **stack**( $a,b$ ) is relevant

by backward application of the action we get a new goal:

$\{\text{holding}(a), \text{clear}(b), \text{on}(b,c)\}$

**stack**( $x,y$ )

Precond:  $\text{holding}(x), \text{clear}(y)$

Effects:  $\sim\text{holding}(x), \sim\text{clear}(y),$   
 $\text{on}(x,y), \text{clear}(x), \text{handempty}$

## Backward planning: algorithm

Backward-search( $O, s_0, g$ )

$\pi \leftarrow$  the empty plan

loop

if  $s_0$  satisfies  $g$  then return  $\pi$

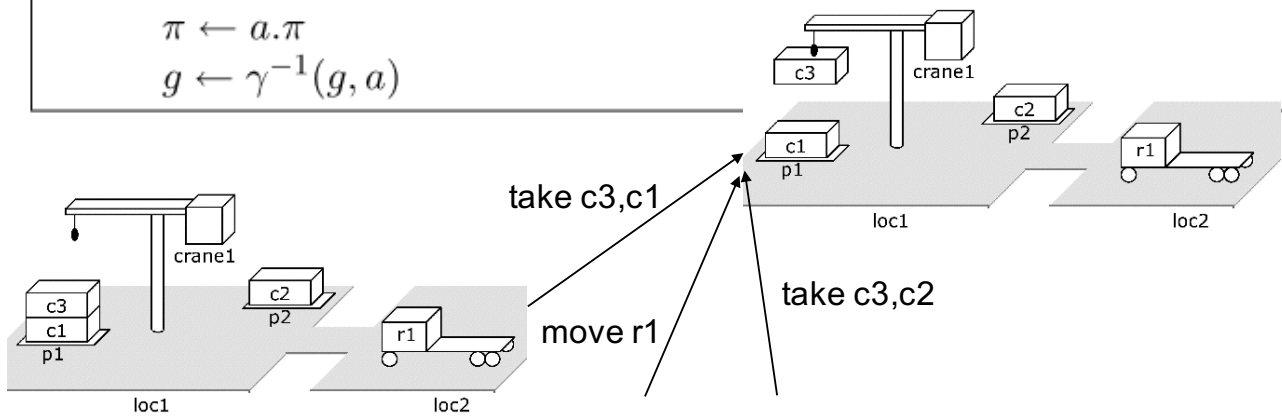
$A \leftarrow \{a \mid a \text{ is a ground instance of an operator in } O$   
and  $\gamma^{-1}(g, a)$  is defined}

if  $A = \emptyset$  then return failure

nondeterministically choose an action  $a \in A$

$\pi \leftarrow a.\pi$

$g \leftarrow \gamma^{-1}(g, a)$



## Backward planning: an example

Goal = {at(r1,loc1),loaded(r1,c3)}

**load(crane1,loc1,c3,r1)**

load( $k, l, c, r$ )

:: crane  $k$  at location  $l$  loads container  $c$  onto robot  $r$

precond: belong( $k, l$ ), holding( $k, c$ ), at( $r, l$ ), unloaded( $r$ )

effects: empty( $k$ ),  $\neg$ holding( $k, c$ ), loaded( $r, c$ ),  $\neg$ unloaded( $r$ )

{at(r1,loc1), belong(crane1,loc1),  
holding(crane1,c3), unloaded(r1)}

**move(r1,loc2,loc1)**

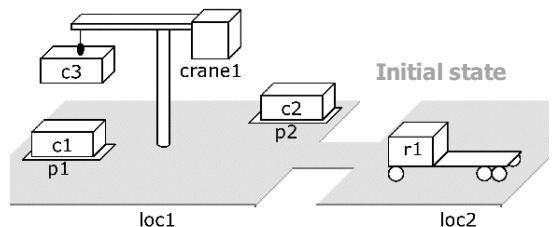
move( $r, l, m$ )

:: robot  $r$  moves from location  $l$  to location  $m$

precond: adjacent( $l, m$ ), at( $r, l$ ),  $\neg$ occupied( $m$ )

effects: at( $r, m$ ), occupied( $m$ ),  $\neg$ occupied( $l$ ),  $\neg$ at( $r, l$ )

{belong(crane1,loc1), holding(crane1,c3),  
unloaded(r1),  
adjacent(loc2,loc1),  
at(r1,loc2),  
 $\neg$ occupied(loc1)}



Backward planning is **sound and complete**.

We can implement a **deterministic** version of the algorithm (via search).

- For completeness we need loop checks.
  - Let  $(g_1, \dots, g_k)$  be a sequence of goals. If  $\exists i < k \ g_i \subseteq g_k$  then we can stop search exploring this branch.

## Branching

- The number of options can be smaller than for the forward planning (less relevant actions for the goal).
- Still, it could be too large.
  - If we want a robot to be at the position loc51 and there are direct connections from states loc1,...,loc50, then we have 50 relevant actions. However, at this stage, the start location is not important!
  - We can instantiate actions only partially (some variables remain free. This is called **lifting**).

```
Lifted-backward-search( $O, s_0, g$ )
 $\pi \leftarrow$  the empty plan
loop
  if  $s_0$  satisfies  $g$  then return  $\pi$ 
   $A \leftarrow \{(o, \theta) \mid o \text{ is a standardization of an operator in } O,$ 
     $\theta \text{ is an mgu for an atom of } g \text{ and an atom of effects } (o),$ 
     $\text{and } \gamma^{-1}(\theta(g), \theta(o)) \text{ is defined}\}$ 
  if  $A = \emptyset$  then return failure
  nondeterministically choose a pair  $(o, \theta) \in A$ 
   $\pi \leftarrow$  the concatenation of  $\theta(o)$  and  $\theta(\pi)$ 
   $g \leftarrow \gamma^{-1}(\theta(g), \theta(o))$ 
```

### Notes:

- standardization = a copy with fresh variables
- mgu = most general unifier
- by using the variables we can decrease the branching factor but the trade off is more complicated loop check

How can we further reduce the search space?

**STRIPS algorithm** reduces the search space of backward planning in the following way:

- **only part of the goal is assumed in each step, namely the preconditions of the last selected action**
  - instead of  $\gamma^{-1}(s, a)$  we can use  $\text{precond}(a)$  as the new goal
  - the rest of the goal must be covered later
  - This makes the algorithm incomplete!
- **If the current state satisfies the preconditions of the selected action then this action is used and never removed later upon backtracking.**

### STRIPS algorithm

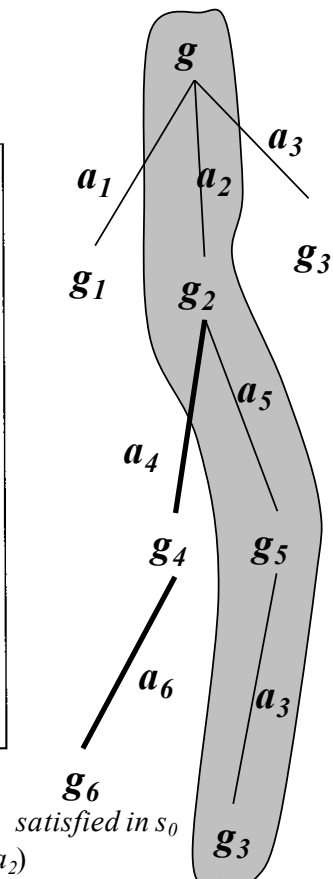
The original STRIPS algorithm is a lifted version of the algorithm below.

```

Ground-STRIPS( $O, s, g$ )
 $\pi \leftarrow$  the empty plan
loop
  if  $s$  satisfies  $g$  then return  $\pi$ 
   $A \leftarrow \{a \mid a \text{ is a ground instance of an operator in } O, \text{ and } a \text{ is relevant for } g\}$ 
  if  $A = \emptyset$  then return failure
  nondeterministically choose any action  $a \in A$ 
   $\pi' \leftarrow \text{Ground-STRIPS}(O, s, \text{precond}(a))$ 
  if  $\pi' = \text{failure}$  then return failure
  ;; if we get here, then  $\pi'$  achieves  $\text{precond}(a)$  from  $s$ 
   $s \leftarrow \gamma(s, \pi')$ 
  ;;  $s$  now satisfies  $\text{precond}(a)$ 
   $s \leftarrow \gamma(s, a)$ 
   $\pi \leftarrow \pi . \pi' . a$ 

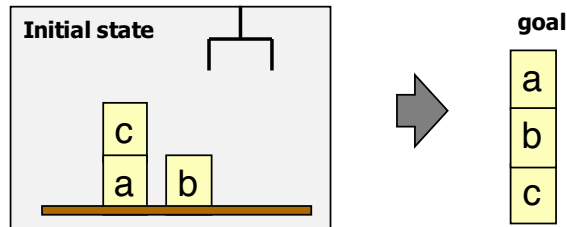
```

$g_2 = (g - \text{effects}(a_2)) \cup \text{precond}(a_2)$   
 $\pi' = \langle a_6, a_4 \rangle$  is a plan for  $\text{precond}(a_2)$   
 $s = \gamma(\gamma(s_0, a_6), a_4)$  is a state satisfying  $\text{precond}(a_2)$



Sussman anomaly is a famous example that causes troubles to the STRIPS algorithm (the algorithm can only find redundant plans).

## Block world



**A plan found by STRIPS** may look like this:

- unstack(c,a), putdown(c), **pickup(a), stack(a,b)**  
*now we satisfied subgoal on(a,b)*
- **unstack(a,b), putdown(a), pickup(b), stack(b,c)**  
*now we satisfied subgoal on(b,c),  
but we need to re-satisfy on(a,b) again*
- pickup(a), stack(a,b) red actions can be deleted

*How to plan for blocks world?*

## Solving Sussman anomaly

### – interleaving plans

- plan-space planning

### – using domain dependent information

- When does a solution plan exist for a blocks world?
  - all blocks from the goal are present in the initial state
  - no block in the goal stays on two other blocks (or on itself)
  - ...
- How to find a solution plan?

Actually, it is easy and very fast!

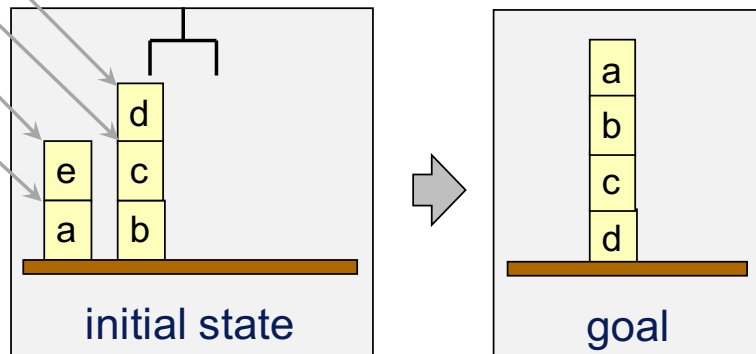
- put all blocks on the table (separately)
- build the requested towers

We can do it even better with additional knowledge!

## When do we need to move block x?

Exactly in one of the following situations:

- s contains **ontable(x)** and g contains **on(x,y)**
- s contains **on(x,y)** and g contains **ontable(x)**
- s contains **on(x,y)** and g contains **on(x,z)** for some  $y \neq z$
- s contains **on(x,y)** and y must be moved

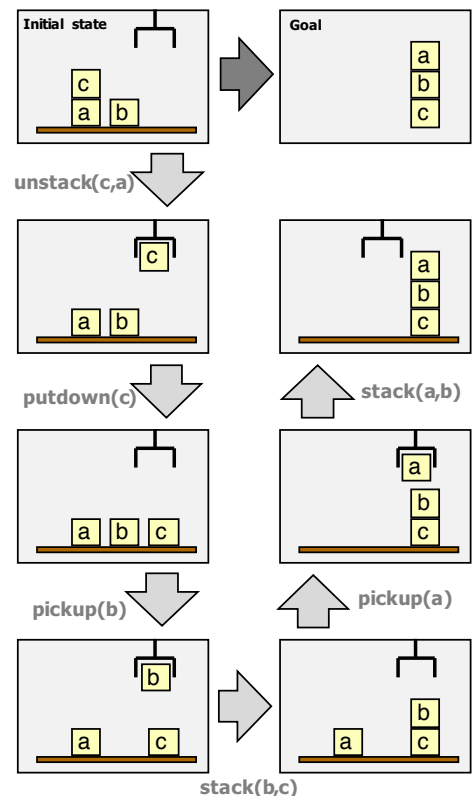


## Fast planning for blocksworld

```

Stack-containers( $O, s_0, g$ ):
  if  $g$  does not satisfy the consistency conditions then
    return failure    ;; the planning problem is unsolvable
   $\pi \leftarrow$  the empty plan
   $s \leftarrow s_0$ 
  loop
    if  $s$  satisfies  $g$  then return  $\pi$ 
    if there are containers  $b$  and  $c$  at the tops of their piles such that
      position( $c, s$ ) is consistent with  $g$  and  $\text{on}(b, c) \in g$ 
    then
      append actions to  $\pi$  that move  $b$  to  $c$ 
       $s \leftarrow$  the result of applying these actions to  $s$ 
      ;; we will never need to move  $b$  again
    else if there is a container  $b$  at the top of its pile
      such that position( $b, s$ ) is inconsistent with  $g$ 
      and there is no  $c$  such that  $\text{on}(b, c) \in g$ 
    then
      append actions to  $\pi$  that move  $b$  to an empty auxiliary pile
       $s \leftarrow$  the result of applying these actions to  $s$ 
      ;; we will never need to move  $b$  again
    else
      nondeterministically choose any container  $c$  such that  $c$  is
      at the top of a pile and position( $c, s$ ) is inconsistent with  $g$ 
      append actions to  $\pi$  that move  $c$  to an empty auxiliary pallet
       $s \leftarrow$  the result of applying these actions to  $s$ 
  
```

Position is consistent with block  $c$  if there is no reason to move  $c$ .



- **Problem Formalisation**
  - models and representations
- **State-space Planning**
  - forward and backward search
- **Plan-space Planning**
  - partial-order planning
- **Control Knowledge in Planning**
  - heuristics
  - control rules



### *Plan space planning: core idea*

The principle of plan space planning is similar to backward planning:

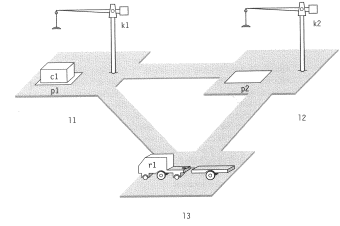
- start from an **„empty“ plan** containing just the description of initial state and goal
- **add other actions** to satisfy not yet covered (open) goals
- if necessary **add other relations** between actions in the plan

Planning is realised as **repairing flaws in a partial plan**

- go from one partial plan to another partial plan until a complete plan is found

Assume a partial plan with the following two actions:

- `take(k1,c1,p1,l1)`
- `load(k1,c1,r1,l1)`



**Possible modifications of the plan:**

- **adding a new action**
  - to apply action **load**, robot `r1` must be at location `l1`
  - action **move**(`r1,l,l1`) moves robot `r1` to location `l1` from some location `l`
- **binding the variables**
  - action **move** is used for the right robot and the right location
- **ordering some actions**
  - the robot must move to the location before the action **load** can be used
  - the order with respect to action **take** is not relevant
- **adding a causal relation**
  - new action is added to move the robot to a given location that is a precondition of another action
  - the causal relation between **move** and **load** ensures that no other action between them moves the robot to another location

**The initial state and the goal** are encoded using two **special actions** in the initial partial plan:

- **Action  $a_0$  represents the initial state** in such a way that atoms from the initial state define effects of the action and there are no preconditions. This action will be before all other actions in the partial plan.
- **Action  $a_\infty$  represents the goal** in a similar way – atoms from the goal define the precondition of that action and there is no effect. This action will be after all other actions.

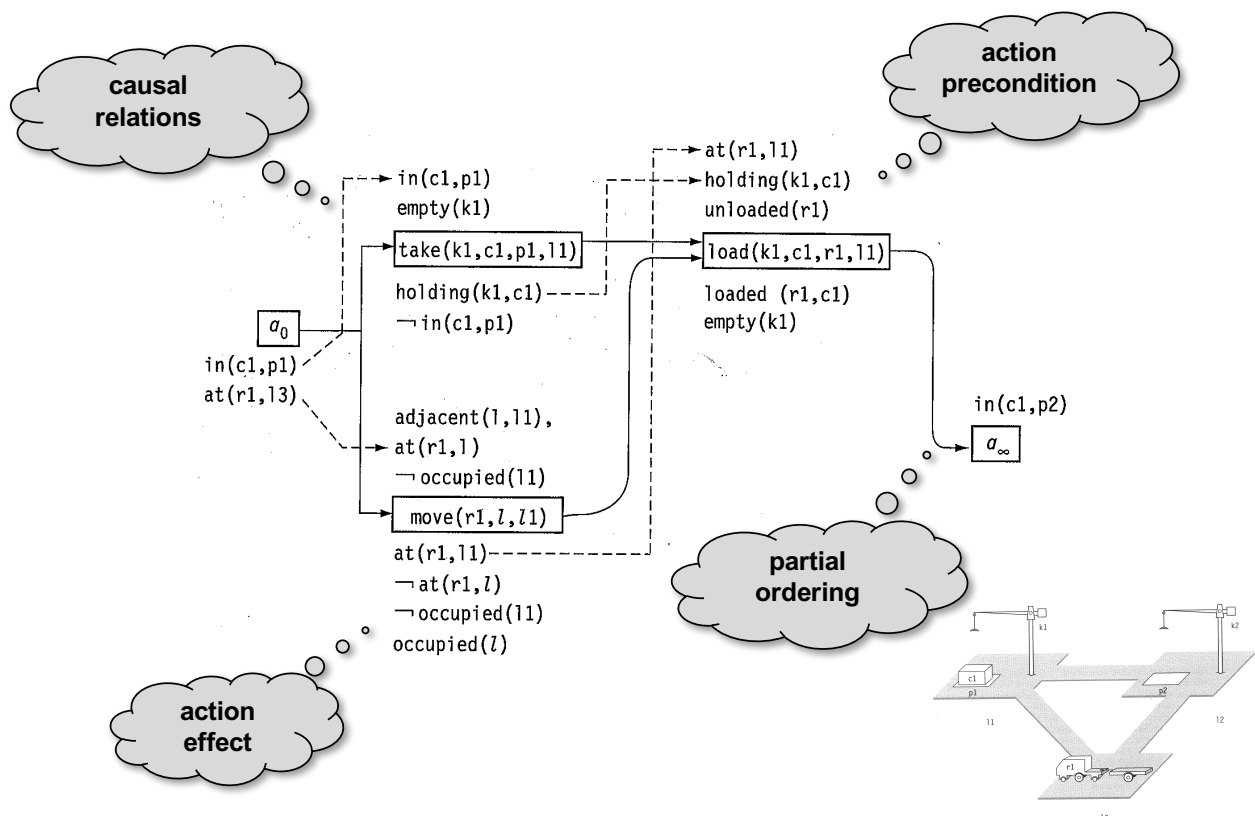
**Planning** is realised by **repairing flaws** in the partial plan.

**The search nodes** correspond to partial plans.

**A partial plan  $\Pi$**  is a tuple  $(A, <, B, L)$ , where

- A is a set of partially instantiated planning operators  $\{a_1, \dots, a_k\}$
- $<$  is a partial order on A ( $a_i < a_j$ )
- B is set of constraints in the form  $x=y$ ,  $x \neq y$  or  $x \in D_i$
- L is a set of causal relations ( $a_i \rightarrow^p a_j$ )
  - $a_i, a_j$  are ordered actions  $a_i < a_j$
  - p is a literal that is effect of  $a_i$  and precondition of  $a_j$
  - B contains relations that bind the corresponding variables in p

*Partial plan: an example*



**Open goal** is an example of a **flaw**.

This is a precondition **p** of some operator **b** in the partial plan such that no action was decided to satisfy this precondition (there is no causal relation  $a_i \rightarrow^p b$ ).

**The open goal p of action b can be resolved by:**

- finding an operator **a** (either present in the partial plan or a new one) that can give **p** (**p** is among the effects of **a** and **a** can be before **b**)
- binding the variables from **p**
- adding a causal relation  $a \rightarrow^p b$

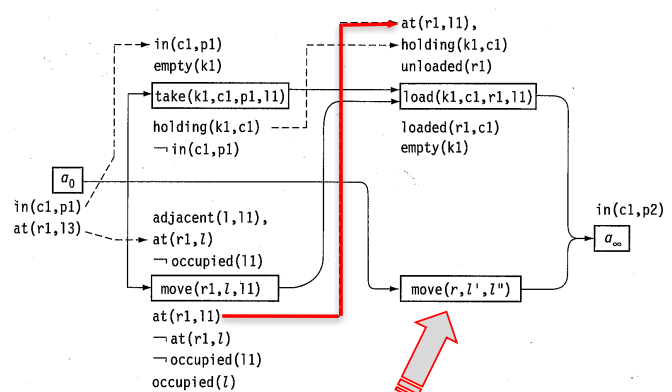
**Threat** is another example of **flaw**.

This is action that can influence existing causal relation.

- Let  $a_i \rightarrow^p a_j$  be a causal relation and action **b** has among its effects a literal unifiable with the negation of **p** and action **b** can be between actions  $a_i$  and  $a_j$ . Then **b** is threat for that causal relation.

We can **remove the threat** by one of the ways:

- ordering **b** before  $a_i$
- ordering **b** after  $a_j$
- binding variables in **b** in such a way that **p** does not bind with the negation of **p**



Partial plan  $\Pi = (A, <, B, L)$  is a **solution plan** for the problem  $P = (\Sigma, s_0, g)$  if:

- partial ordering  $<$  and constraints  $B$  are globally consistent
  - there are no cycles in the partial ordering
  - we can assign variables in such a way that constraints from  $B$  hold
- Any linearly ordered sequence of fully instantiated actions from  $A$  satisfying  $<$  and  $B$  goes from  $s_0$  to a state satisfying  $g$ .

Hmm, but this definition **does not say how** to verify that a partial plan is a solution plan!

*Solution plan – a constructive view*

**How to efficiently verify that a partial plan is a solution plan?**

**Claim:**

Partial plan  $\Pi = (A, <, B, L)$  is a solution plan if:

- there are no flaws (no open goals and no threats)
- partial ordering  $<$  and constraints  $B$  are globally consistent

**Proof** by induction using the plan length

- $\{a_0, a_1, a_\infty\}$  is a solution plan
- for more actions take one of the possible first actions and join it with action  $a_0$

**PSP = Plan-Space Planning**

```

PSP( $\pi$ )
   $flaws \leftarrow \text{OpenGoals}(\pi) \cup \text{Threats}(\pi)$ 
  if  $flaws = \emptyset$  then return( $\pi$ )
  select any flaw  $\phi \in flaws$ 
   $resolvers \leftarrow \text{Resolve}(\phi, \pi)$ 
  if  $resolvers = \emptyset$  then return(failure)
  nondeterministically choose a resolver  $\rho \in resolvers$ 
   $\pi' \leftarrow \text{Refine}(\rho, \pi)$ 
  return(PSP( $\pi'$ ))
end

```

*Notes:*

- The selection of flaw is deterministic (all flaws must be resolved).
- The resolver is selected non-deterministically (search in case of failure).

*PSP – some details*

**Open goals** can be maintained in an **agenda** of action preconditions without causal relations. Adding a causal relation for **p** removes **p** from the agenda.

**All threats** can be found in time  $O(n^3)$  by verifying triples of actions or threats can be maintained incrementally: after adding a new action, check causal relations influenced ( $O(n^2)$ ), after adding a causal relation find its threats ( $O(n)$ ).

Open goals and threats are resolved only by **consistent refinements** of the partial plan.

- consistent ordering can be detected by finding cycles or by maintaining a transitive closure of  $<$
- consistency of constraints in B
  - If there is no negation then we can use arc consistency.
  - In case of negation, the problem of checking global consistency is NP-complete.

## Algorithm PSP is **complete and sound**.

### – **soundness**

- If the algorithm finishes, it returns a consistent plan with no flaws so it is a solution plan.

### – **completeness**

- If there is a solution plan then the algorithm has the option to select the right actions to the partial plan.

## Be careful about the **deterministic implementation!**

### – **The search space is not finite!**

- A complete deterministic procedure must guarantee that it eventually finds a solution plan of any length – **iterative deepening** can be applied.

## Algorithm PoP

## PoP is a popular instance of algorithm PSP.

```

PoP( $\pi$ , agenda)      ;; where  $\pi = (A, <, B, L)$ 
if agenda =  $\emptyset$  then return( $\pi$ )
select any pair ( $a_i, p$ ) in and remove it from agenda
relevant  $\leftarrow$  Providers( $p, \pi$ )
if relevant =  $\emptyset$  then return(failure)
nondeterministically choose an action  $a_i \in$  relevant
 $L \leftarrow L \cup \{(a_i \xrightarrow{p} a_j)\}$ 
update B with the binding constraints of this causal link
if  $a_i$  is a new action in A then do:
  update A with  $a_i$ 
  update  $<$  with ( $a_i < a_j$ ), ( $a_0 < a_i < a_\infty$ )
  update agenda with all preconditions of  $a_i$ 
for each threat on  $\langle a_i \xrightarrow{p} a_j \rangle$  or due to  $a_i$  do:
  resolvers  $\leftarrow$  set of resolvers for this threat
  if resolvers =  $\emptyset$  then return(failure)
  nondeterministically choose a resolver in resolvers
  add that resolver to  $<$  or to B
return(PoP( $\pi$ , agenda))
end
  
```

- **Agenda** is a set of pairs (**a**,**p**), where **p** is an open precondition of action **a**.
- **First find an action  $a_i$  to cover some  $p$  from the agenda.**
- **At the second stage resolve all threats that appeared by adding action  $a_i$  or from a causal relation with  $a_i$ .**

## Initial state:

- $\text{At}(\text{Home})$ ,  $\text{Sells}(\text{OBI}, \text{Drill})$ ,  $\text{Sells}(\text{Tesco}, \text{Milk})$ ,  $\text{Sells}(\text{Tesco}, \text{Banana})$
- so action **Start** is defined as:
  - Precond: none
  - Effects:  $\text{At}(\text{Home})$ ,  $\text{Sells}(\text{OBI}, \text{Drill})$ ,  $\text{Sells}(\text{Tesco}, \text{Milk})$ ,  $\text{Sells}(\text{Tesco}, \text{Banana})$

## Goal:

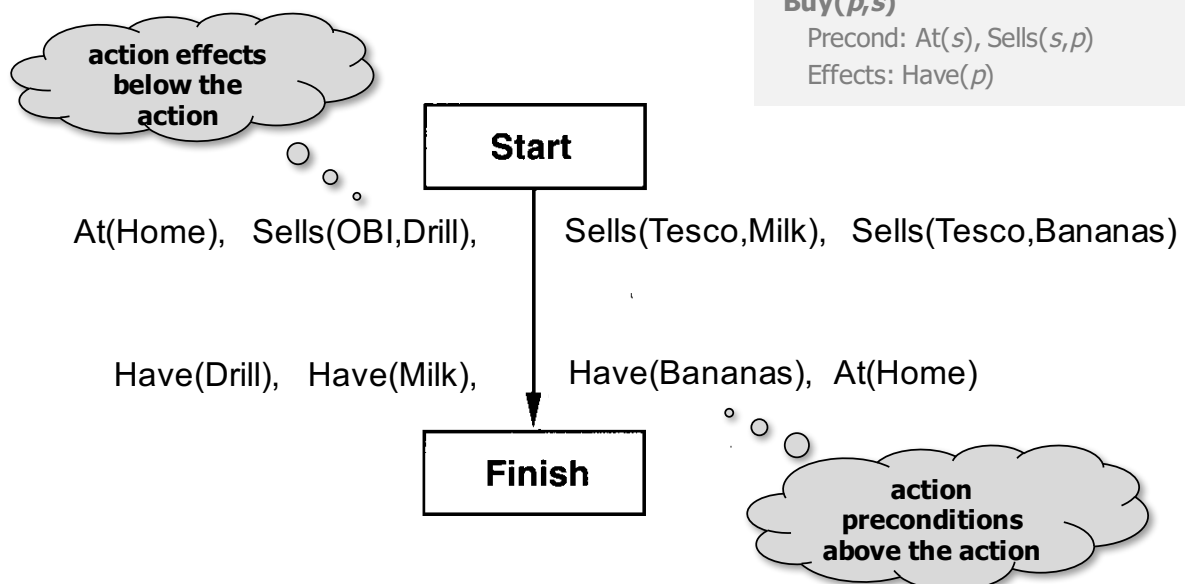
- $\text{Have}(\text{Drill})$ ,  $\text{Have}(\text{Milk})$ ,  $\text{Have}(\text{Banana})$ ,  $\text{At}(\text{Home})$
- so action **Finish** is defined as:
  - Precond:  $\text{Have}(\text{Drill})$ ,  $\text{Have}(\text{Milk})$ ,  $\text{Have}(\text{Banana})$ ,  $\text{At}(\text{Home})$
  - Effects: none

The following two **operators** are available:

- **Go( $l, m$ )** ;; go from location  $l$  to  $m$ 
  - Precond:  $\text{At}(l)$
  - Effects:  $\text{At}(m)$ ,  $\neg \text{At}(l)$
- **Buy( $p, s$ )** ;; buy  $p$  at location  $s$ 
  - Precond:  $\text{At}(s)$ ,  $\text{Sells}(s, p)$
  - Effects:  $\text{Have}(p)$



## The initial (empty) plan



## Plan-space planning: a running example

There is only one way to satisfy the **open goals Have**, and this is via **actions Buy** (no threats added).

### Operators

#### Go(*l*,*m*)

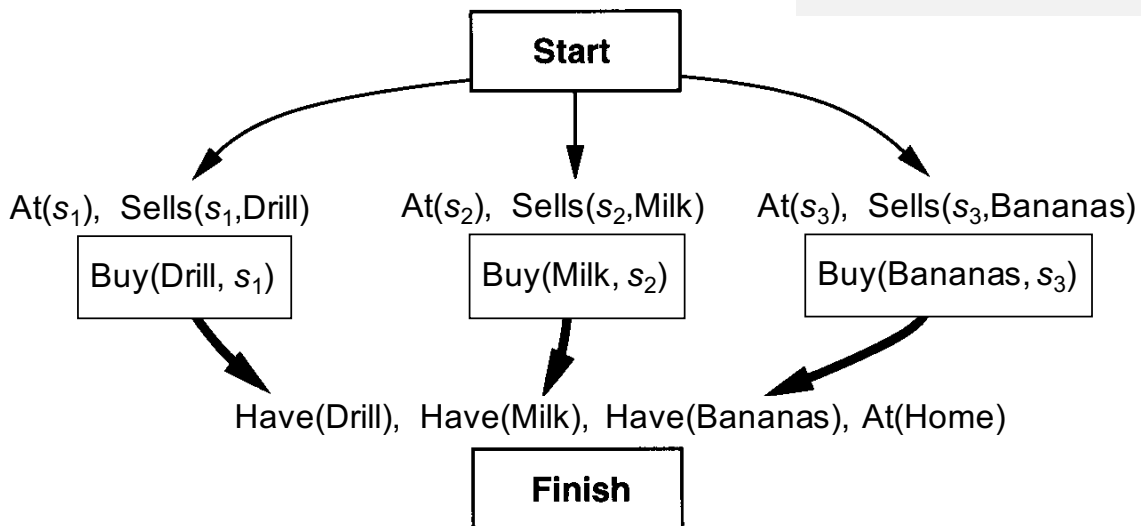
Precond: At(*l*)

Effects: At(*m*),  $\neg$ At(*l*)

#### Buy(*p*,*s*)

Precond: At(*s*), Sells(*s*,*p*)

Effects: Have(*p*)



## Plan-space planning: a running example

There is again a single way to satisfy preconditions **Sells** and this is substituting the right **constants**.

### Operators

#### Go(*l*,*m*)

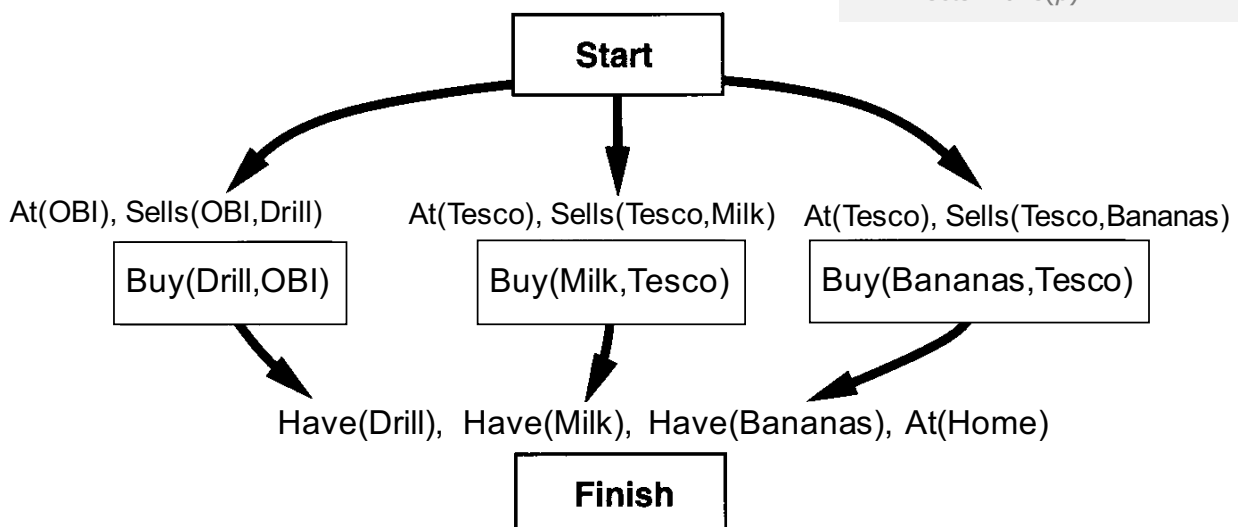
Precond: At(*l*)

Effects: At(*m*),  $\neg$ At(*l*)

#### Buy(*p*,*s*)

Precond: At(*s*), Sells(*s*,*p*)

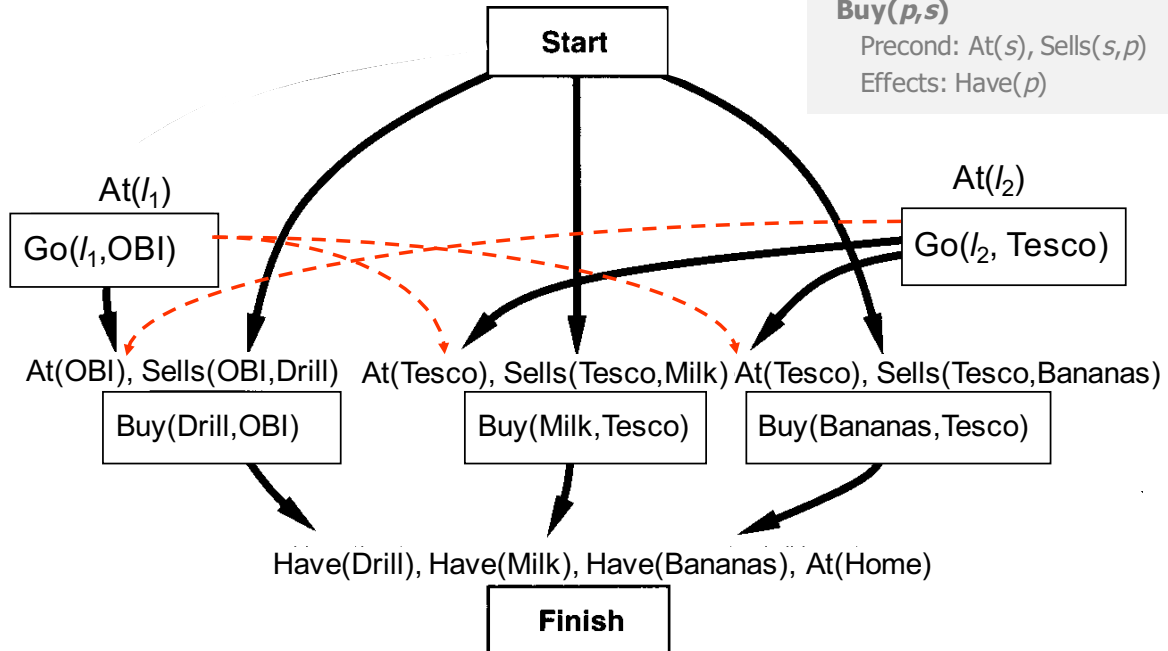
Effects: Have(*p*)



## Plan-space planning: a running example

The only way to **satisfy open goals** is by adding actions **Go**.

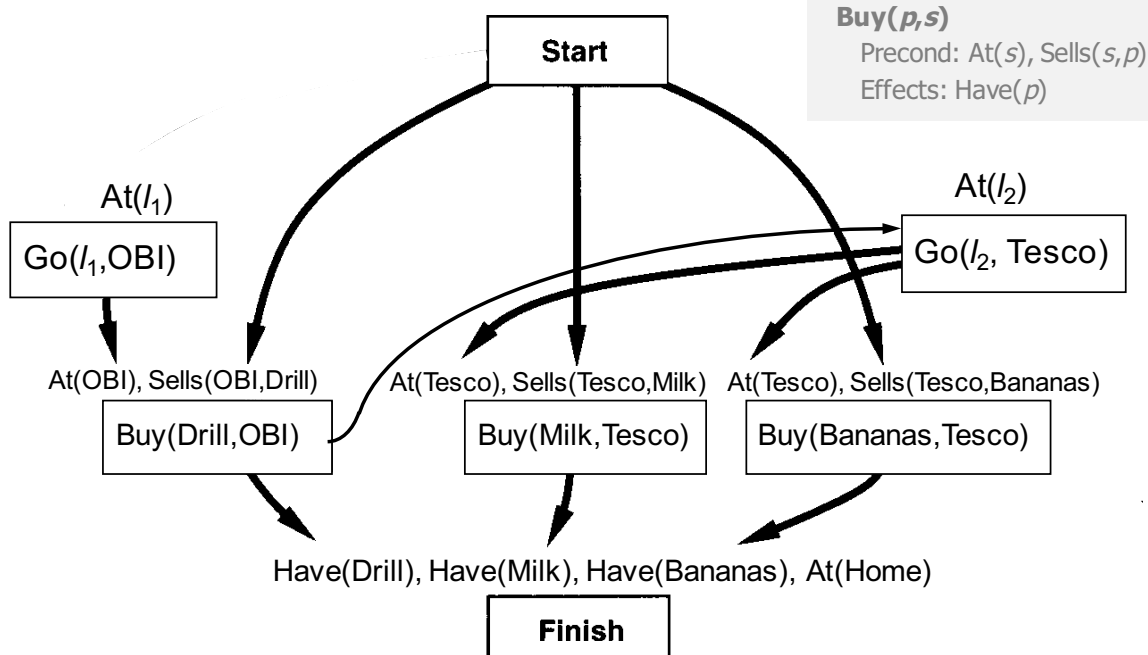
- There are new threats!



## Plan-space planning: a running example

One **threat** can be solved by ordering  $Buy(Drill)$  before  $Go(Tesco)$

- This solves all the threats!



## Plan-space planning: a running example

**Open goal  $At(I_1)$**  can be satisfied by **assignment  $I_1=Home$**  taken from the action **Start**.

### Operators

#### **Go( $l, m$ )**

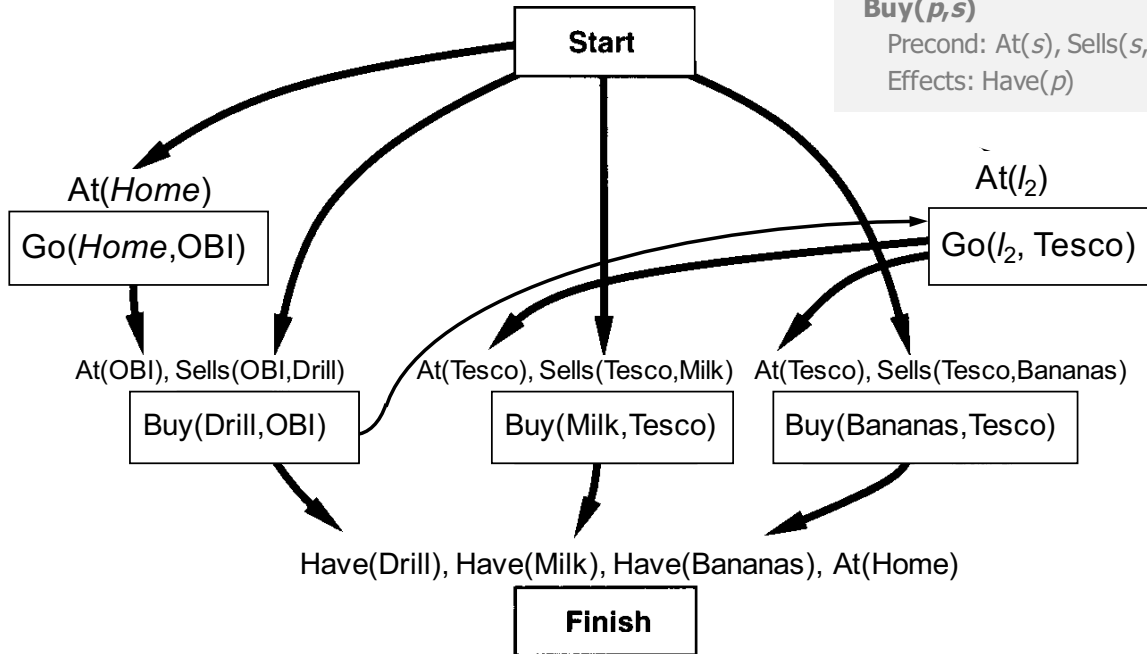
Precond:  $At(l)$

Effects:  $At(m), \neg At(l)$

#### **Buy( $p, s$ )**

Precond:  $At(s), Sells(s, p)$

Effects:  $Have(p)$



## Plan-space planning: a running example

**Open goal  $At(I_2)$**  can be satisfied by **assignment  $I_2=OBI$**  from action **Go( $Home, OBI$ )**

### Operators

#### **Go( $l, m$ )**

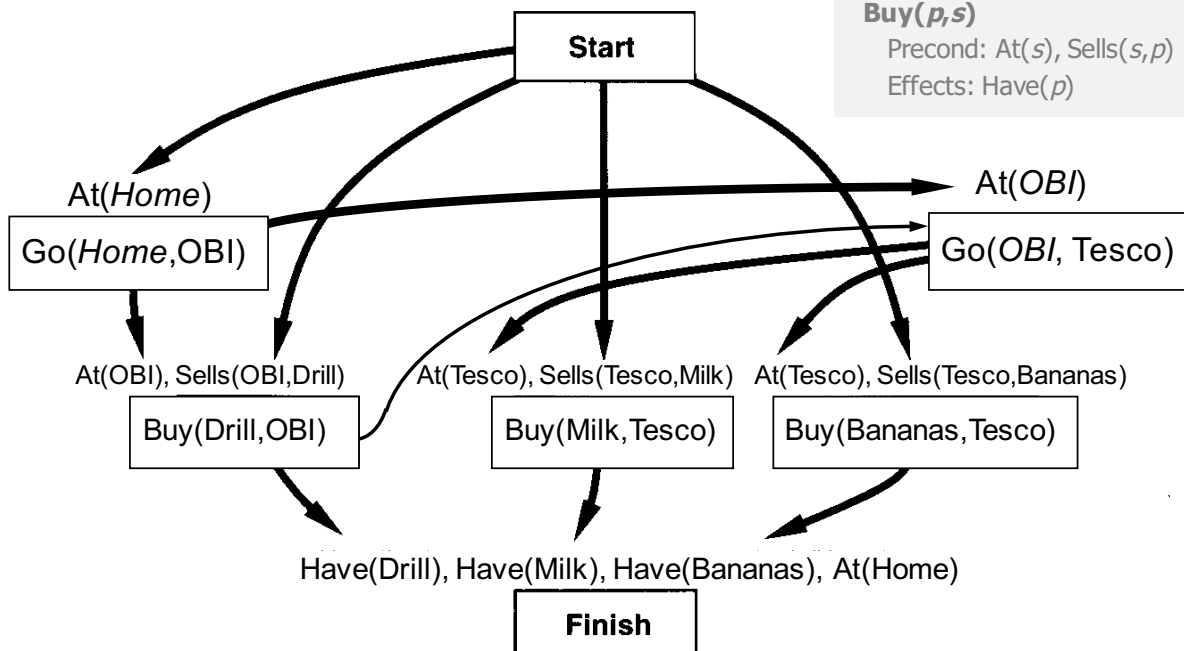
Precond:  $At(l)$

Effects:  $At(m), \neg At(l)$

#### **Buy( $p, s$ )**

Precond:  $At(s), Sells(s, p)$

Effects:  $Have(p)$



## Plan-space planning: a running example

**Open goal**  $At(Home)$  from Finish is satisfied by **action** Go

– new threats appear

### Operators

#### $Go(l, m)$

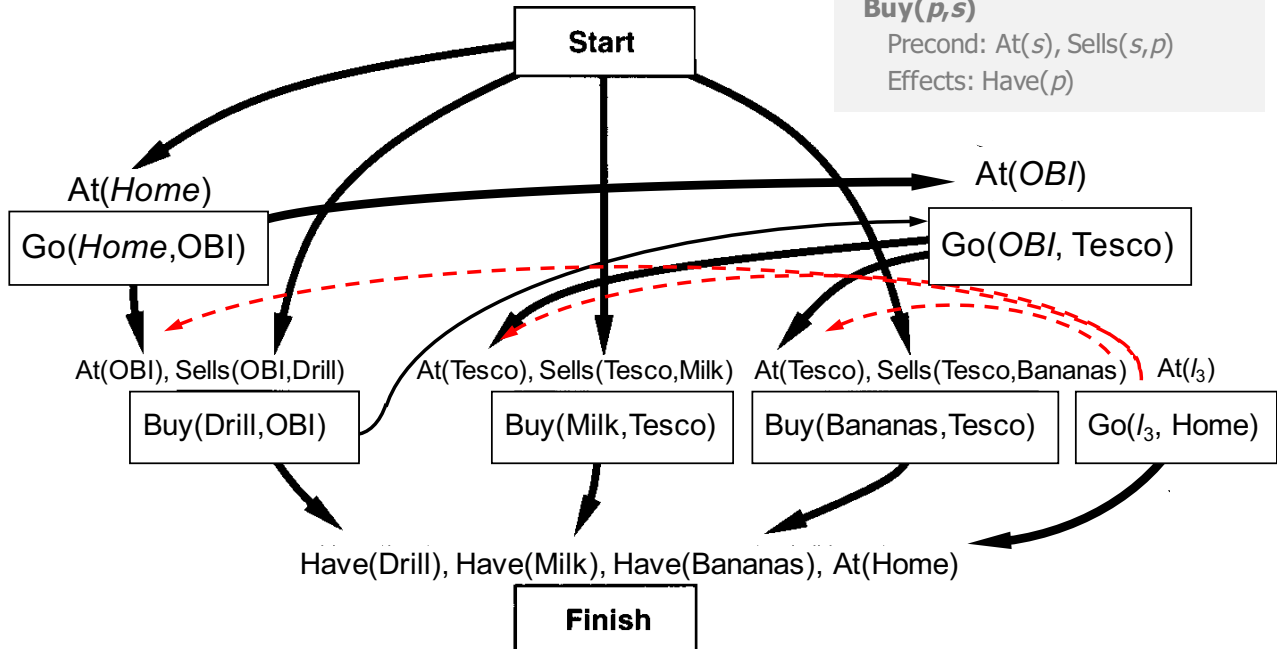
Precond:  $At(l)$

Effects:  $At(m), \neg At(l)$

#### $Buy(p, s)$

Precond:  $At(s), Sells(s, p)$

Effects:  $Have(p)$



## Plan-space planning: a running example

**Threats** for  $At(Tesco)$  are removed by **ordering**  $Go(Home)$  after both actions **Buy**

### Operators

#### $Go(l, m)$

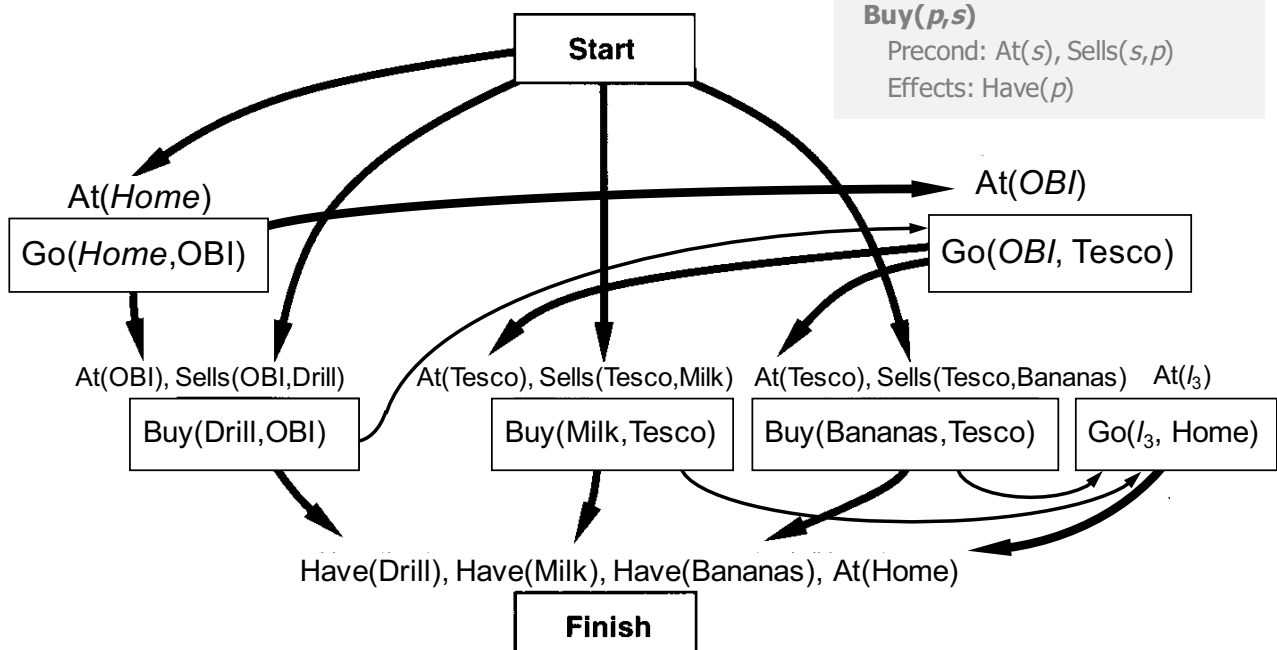
Precond:  $At(l)$

Effects:  $At(m), \neg At(l)$

#### $Buy(p, s)$

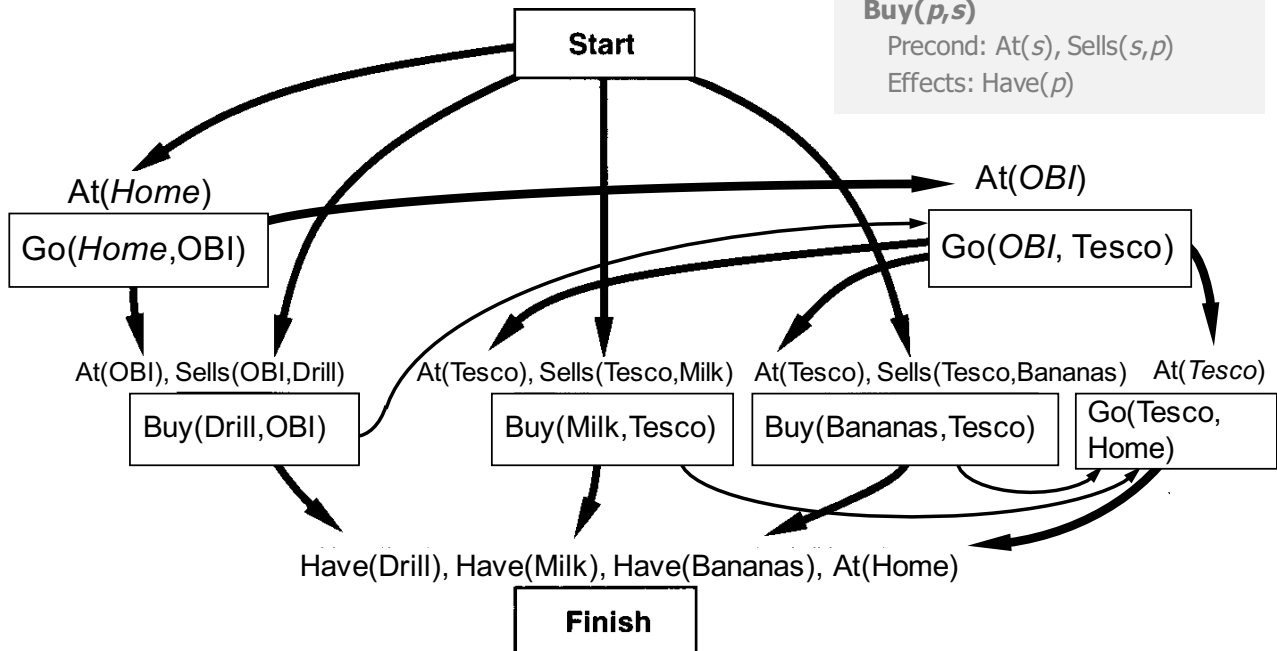
Precond:  $At(s), Sells(s, p)$

Effects:  $Have(p)$



## Plan-space planning: a running example

**Open goal**  $At(I_3)$  is satisfied by **assignment**  $I_3 = \text{Tesco}$  from action  $\text{Go}(\text{OBI}, \text{Tesco})$ .



## Comparison

	State space planning	Plan space planning
search space	finite	infinite
search nodes	simple (world states)	complex (partial plans)
world states	explicit	not used
partial plan	action selection and ordering done together	action selection and ordering separated
plan structure	linear	causal relations

**State space planning is much faster** today thanks to heuristics based on state evaluation.

However, **plan space planning**:

- makes more **flexible plans** thanks to partial order
- supports **further extensions** such as adding explicit time and resources

- **Problem Formalisation**
  - models and representations
- **State-space Planning**
  - forward and backward search
- **Plan-space Planning**
  - partial-order planning
- **Control Knowledge in Planning**
  - heuristics
  - control rules



## Heuristics

**Heuristics** are used to select next search node to be explored (recall, that we described the planning algorithms using non-determinism).

- Note: If we know, which node to select to get a solution, then we use **oracle**. With oracle we will find the solution deterministically.

Naturally, we prefer the heuristic to be as **close** as possible to **oracle** while being **computed efficiently**.

A typical way to obtain (admissible) heuristics is via solving a **relaxed problem** (some problem constraints are relaxed – not assumed).

- solve the relaxed problem for the successor nodes
- select the node with the best solution of the relaxed problem

For optimisation problems the heuristic  $h(u)$  estimates the real cost  $h^*(u)$  of the best solution reachable via node  $u$ .

- the heuristic is **admissible**, if  $h(u) \leq h^*(u)$  (for minimization)
- the search algorithms using admissible heuristics are optimal

Heuristic estimates the **number of actions** to reach a goal state from a given state or to reach a given predicate or a set of predicates.

Based on solving a **“relaxed” problem**:

- assume only positive effects
- assume that different atoms can be reached independently

**Zero attempt:**

- $\Delta_0(s,p) = 0$  if  $p \in s$
- $\Delta_0(s,g) = 0$  if  $g \subseteq s$
- $\Delta_0(s,p) = \infty$  if  $p \notin s$  and  $\forall a \in A, p \notin \text{effects}^+(a)$
- $\Delta_0(s,p) = \min_a \{1 + \Delta_0(s, \text{precond}(a)) \mid p \in \text{effects}^+(a)\}$
- $\Delta_0(s,g) = \sum_{p \in g} \Delta_0(s,p)$

This heuristic is **not admissible** (for optimal planning) because it does not provide a lower bound for the plan length!

```

Delta(s)
  for each p do: if p ∈ s then Δ0(s,p) ← 0, else Δ0(s,p) ← ∞
  U ← s
  iterate
    for each a such that precond(a) ⊆ U do
      U ← U ∪ effects+(a)
      for each p ∈ effects+(a) do
        Δ0(s,p) ← min{Δ0(s,p), 1 + ∑q ∈ precond(a) Δ0(s,q)}
      until no change occurs in the above updates
  end
  
```

**A first attempt to admissible heuristic**

- ...
- $\Delta_1(s,g) = \max\{\Delta_0(s,p) \mid p \in g\}$
- If the heuristic value is greater than the best so-far solution then we can cut-off the search branch.
- Based on experiments, heuristic  $\Delta_1$  is less informed than  $\Delta_0$ .

**A second attempt to admissible heuristic**

Let us try to explore reachability of pairs of atoms together.

- ...
- $\Delta_2(s,p) = \min_a \{1 + \Delta_2(s, \text{precond}(a)) \mid p \in \text{effects}^+(a)\}$
- $\Delta_2(s,\{p,q\}) = \min\{$   
 $\min_a \{1 + \Delta_2(s, \text{precond}(a)) \mid \{p,q\} \subseteq \text{effects}^+(a)\},$   
 $\min_a \{1 + \Delta_2(s, \{q\} \cup \text{precond}(a)) \mid p \in \text{effects}^+(a)\},$   
 $\min_a \{1 + \Delta_2(s, \{p\} \cup \text{precond}(a)) \mid q \in \text{effects}^+(a)\}\}$
- $\Delta_2(s,g) = \max_{p,q} \{\Delta_2(s,\{p,q\}) \mid \{p,q\} \subseteq g\}$

We can generalise the above idea to larger sets of atoms, but for  $k > 2$  this heuristic is computationally expensive.

## Forward planning

- Prefer the action leading to a state with smaller heuristic distance to a goal.
- Heuristic is computed in every search step.

```

Heuristic-forward-search( $\pi, s, g, A$ )
  if  $s$  satisfies  $g$  then return  $\pi$ 
  options  $\leftarrow \{a \in A \mid a \text{ applicable to } s\}$ 
  for each  $a \in \text{options}$  do Delta( $\gamma(s, a)$ )
  while options  $\neq \emptyset$  do
     $a \leftarrow \text{argmin}\{\Delta_0(\gamma(s, a), g) \mid a \in \text{options}\}$ 
    options  $\leftarrow \text{options} - \{a\}$ 
     $\pi' \leftarrow \text{Heuristic-forward-search}(\pi, a, \gamma(s, a), g, A)$ 
    if  $\pi' \neq \text{failure}$  then return( $\pi'$ )
  return(failure)
end
    
```

## Backward planning

- First, compute the heuristic distance from the initial state  $s_0$  to all atoms:  $\Delta(s_0, p)$ 
  - can be done incrementally
- Prefer the action whose regression set is heuristically closer to the initial state.

```

Backward-search( $\pi, s_0, g, A$ )
  if  $s_0$  satisfies  $g$  then return( $\pi$ )
  options  $\leftarrow \{a \in A \mid a \text{ relevant for } g\}$ 
  while options  $\neq \emptyset$  do
     $a \leftarrow \text{argmin}\{\Delta_0(s_0, \gamma^{-1}(g, a)) \mid a \in \text{options}\}$ 
    options  $\leftarrow \text{options} - \{a\}$ 
     $\pi' \leftarrow \text{Backward-search}(a, \pi, s_0, \gamma^{-1}(g, a), A)$ 
    if  $\pi' \neq \text{failure}$  then return( $\pi'$ )
  return failure
end
    
```

## Plan-space heuristics

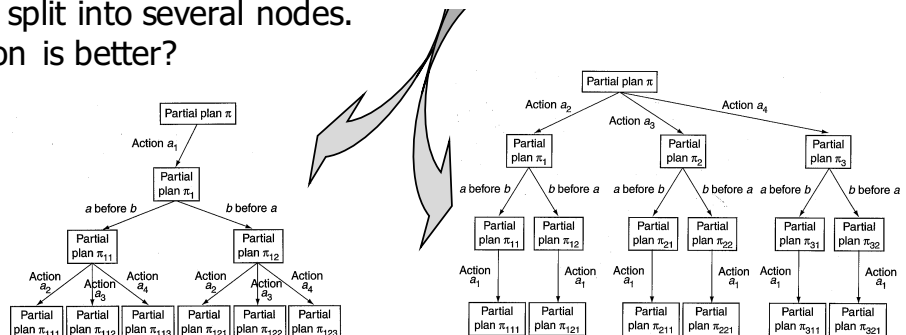
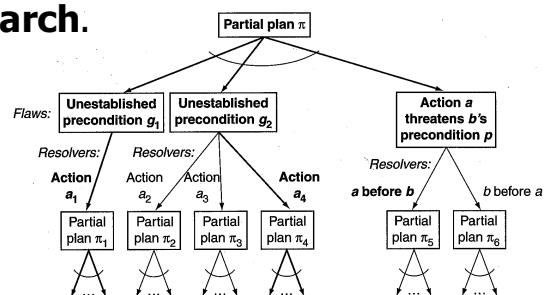
Plan-space planning is based on **AND-OR search**.

There are two types of choices:

- the choice of flaw (AND node)
- the choice of resolver (OR node)

## Flaw-selection heuristic

- This is a form of **serialization of the AND/OR tree**, in particular the AND node is split into several nodes.
- Which serialization is better?



- Better serialization leads to a smaller number of nodes in the graph.
- **FAF (fewest alternatives first) heuristic**
  - first repair the flaws with fewer ways for repair

### Which resolver for a flaw should be tried first?

Let  $\{\pi_1, \dots, \pi_m\}$  be partial plans obtained by applying different flaw resolvers and  $g_\pi$  be a set of open goals in  $\pi$ .

- **Zero attempt**

prefer a partial plan with fewer open goals

$$\Leftrightarrow \eta_0(\pi) = |g_\pi|$$

- However, this does not really estimate the size of the plan.

- **Next attempt**

Generate an AND-OR graph for  $\pi$  till given depth  $k$  and count the number of new actions and the number of open goals not in  $s_0$

$$\Leftrightarrow \eta_k(\pi)$$

- This is **too computationally expensive**.

- **One more improvement**

Construct a planning graph (once) for the original goal. Then find an open goal  $p$  in  $\pi$ , that was added last to the graph and on the path from  $s_0$  to  $p$  count the number of actions that are not in  $\pi$

$$\Leftrightarrow \eta(\pi)$$

## Pruning

Heuristics guide the planner towards a goal state by ordering alternative plans. They do not solve the problem with the **large number of alternatives**.

Can we **detect and prune bad alternatives**?

### Example (blockworld)

- If a block is placed correctly (consistent with the goal) then any action that moves that block just enlarges the plan.
- If a block is on a wrong place and there is an action that moves it to the correct place then any action that moves the block elsewhere just enlarges the plan.

Domain dependent information can prune the search space, but the open question is how to express such information for a general planning algorithm.

- **control rules**

We need a formalism to express relations between the current world state and future states.

## Simple temporal logic

- extension of first-order logic by **modal operators**
  - $\phi_1 \cup \phi_2$  (until)       $\phi_1$  is true in all states until the first state (if any) in which  $\phi_2$  is true
  - $\Box \phi$  (always)       $\phi$  is true now and in all future states
  - $\Diamond \phi$  (eventually)       $\phi$  is true now or in any future state
  - $\bigcirc \phi$  (next)       $\phi$  is true in the next state
  - $\text{GOAL}(\phi)$        $\phi$  (no modal operators) is true in the goal state
- $\phi$  is a logical formula expressing relations between the objects of the world (it can include modal operators)

## Semantics of modal operators

The **interpretation** of modal formula involves not just the current state but we need to work with a triple **(S, s<sub>i</sub>, g)**:

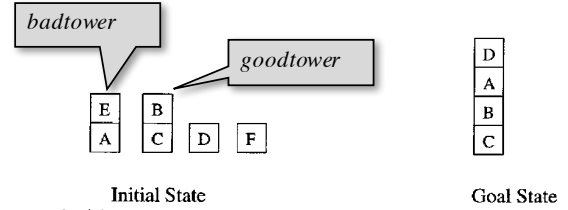
- $S = \langle s_0, s_1, \dots \rangle$  is an infinite sequence of states
- $s_i \in S$  is the current state
- $g$  is a goal formula

Plan  $\pi = \langle a_1, a_2, \dots, a_n \rangle$  gives a finite sequence of states  $S_\pi = \langle s_0, s_1, \dots, s_n \rangle$  where  $s_{i+1} = \gamma(s_i, a_{i+1})$ , that can be made infinite  $\langle s_0, s_1, \dots, s_{n-1}, s_n, s_n, s_n, \dots \rangle$

$(S, s_i, g) \models \phi$  is defined as follows:

- $(S, s_i, g) \models \phi$       iff  $s_i \models \phi$  for atom  $\phi$
- $(S, s_i, g) \models \phi_1 \wedge \phi_2$       iff  $(S, s_i, g) \models \phi_1$  and  $(S, s_i, g) \models \phi_2$
- ...
- $(S, s_i, g) \models \phi_1 \cup \phi_2$       iff there exists  $j \geq i$  st.  $(S, s_j, g) \models \phi_2$  and for each  $k: i \leq k < j$   $(S, s_k, g) \models \phi_1$
- $(S, s_i, g) \models \Box \phi$       iff  $(S, s_j, g) \models \phi$  for each  $j \geq i$
- $(S, s_i, g) \models \Diamond \phi$       iff  $(S, s_j, g) \models \phi$  for some  $j \geq i$
- $(S, s_i, g) \models \bigcirc \phi$       iff  $(S, s_{i+1}, g) \models \phi$
- $(S, s_i, g) \models \text{GOAL}(\phi)$       iff  $\phi \in g$

*Goodtower* is a tower such that no block needs to be moved.  
*Badtower* is a tower that is not good.



$$goodtower(x) \triangleq clear(x) \wedge \neg GOAL(holding(x)) \wedge goodtowerbelow(x)$$

$$goodtowerbelow(x) \triangleq (ontable(x) \wedge \neg \exists[y:GOAL(on(x, y))]) \vee \exists[y:on(x, y)] \neg GOAL(ontable(y)) \wedge \neg GOAL(holding(y)) \wedge \neg GOAL(clear(y)) \wedge \forall[z:GOAL(on(x, z))] z = y \wedge \forall[z:GOAL(on(z, y))] z = x \wedge goodtowerbelow(y)$$

$$badtower(x) \triangleq clear(x) \wedge \neg goodtower(x)$$

**Control rule:**

$$\square \left( \begin{aligned} & \forall[x:clear(x)] goodtower(x) \Rightarrow \bigcirc(clear(x) \vee \exists[y:on(y, x)] goodtower(y)) \\ & \wedge badtower(x) \Rightarrow \bigcirc(\neg \exists[y:on(y, x)]) \\ & \wedge (ontable(x) \wedge \exists[y:GOAL(on(x, y))] \neg goodtower(y)) \\ & \Rightarrow \bigcirc(\neg holding(x)) \end{aligned} \right)$$

goodtower remains goodtower

do not put anything on badtower

do not take a block from a table until you can put it on a goodtower

## Progression

To use control rules in planning we need to express how the formula changes when we **go from state  $s_i$  to state  $s_{i+1}$** .

- We look for a formula  $progr(\phi, s_i)$  that is true in  $s_{i+1}$ , if  $\phi$  is true in state  $s_i$
- $\phi$  does not contain any modal operator
  - $progr(\phi, s_i) = \text{true}$  if  $s_i \models \phi$
  - $= \text{false}$  if  $s_i \not\models \phi$  does not hold
- $\phi$  with logical connectives
  - $progr(\phi_1 \wedge \phi_2, s_i) = progr(\phi_1, s_i) \wedge progr(\phi_2, s_i)$
  - $progr(\neg \phi, s_i) = \neg progr(\phi, s_i)$
- $\phi$  with quantifiers (no function symbols, just  $k$  constants  $c_j$ )
  - $progr(\forall x \phi, s_i) = progr(\phi\{x/c_1\}, s_i) \wedge \dots \wedge progr(\phi\{x/c_k\}, s_i)$
  - $progr(\exists x \phi, s_i) = progr(\phi\{x/c_1\}, s_i) \vee \dots \vee progr(\phi\{x/c_k\}, s_i)$
- $\phi$  with modal operators
  - $progr(\phi_1 \cup \phi_2, s_i) = ((\phi_1 \cup \phi_2) \wedge progr(\phi_1, s_i)) \vee progr(\phi_2, s_i)$
  - $progr(\Box \phi, s_i) = (\Box \phi) \wedge progr(\phi, s_i)$
  - $progr(\Diamond \phi, s_i) = (\Diamond \phi) \vee progr(\phi, s_i)$
  - $progr(\bigcirc \phi, s_i) = \phi$

**Technical notes:**

- $progress(\phi, s_i)$  is obtained from  $progr(\phi, s_i)$  by cleaning ( $\text{true} \wedge d \rightarrow d$ ,  $\neg \text{true} \rightarrow \text{false}$ , ...)
- Can be extended to a sequence of states  $\langle s_0, \dots, s_n \rangle$ 

$$progress(\phi, \langle s_0, \dots, s_n \rangle) = \begin{cases} \phi & \text{if } n = 0 \\ progress(progress(\phi, \langle s_0, \dots, s_{n-1} \rangle), s_n) & \text{otherwise} \end{cases}$$

$(S, s_i, g) \vdash \phi$  iff  $(S, s_{i+1}, g) \vdash \text{progress}(\phi, s_i)$ .

- i.e. progress behaves as we need

$(S, s_0, g) \vdash \phi$  then for any prefix  $S' = \langle s_0, \dots, s_i \rangle$  of  $S$  it holds  $\text{progress}(\phi, S') \neq \text{false}$ .

- If the control rule is satisfied then progress is not false

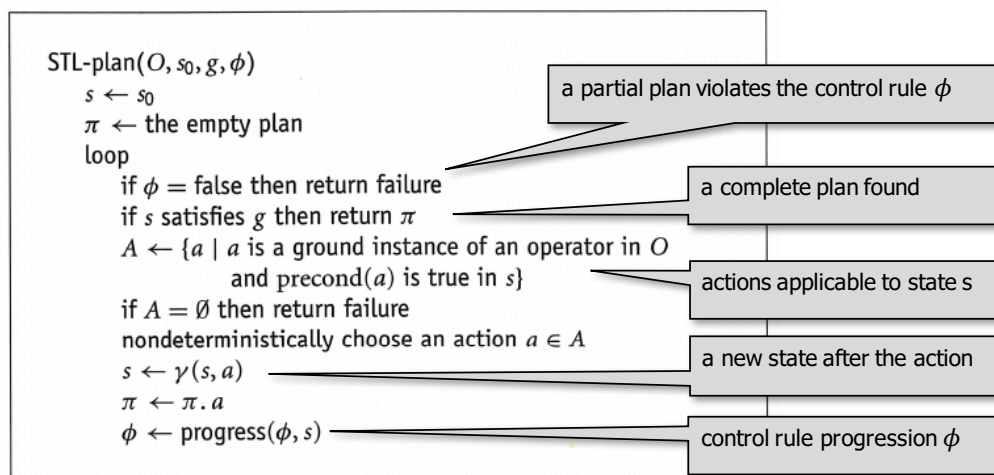
If plan  $\pi$  is applicable to  $s_0$  and  $\text{progress}(\phi, S_\pi) = \text{false}$ , then there is no extension  $S'$  of  $S_\pi$  st.  $(S', s_0, g) \vdash \phi$

- If progress is false then the control rule cannot be satisfied

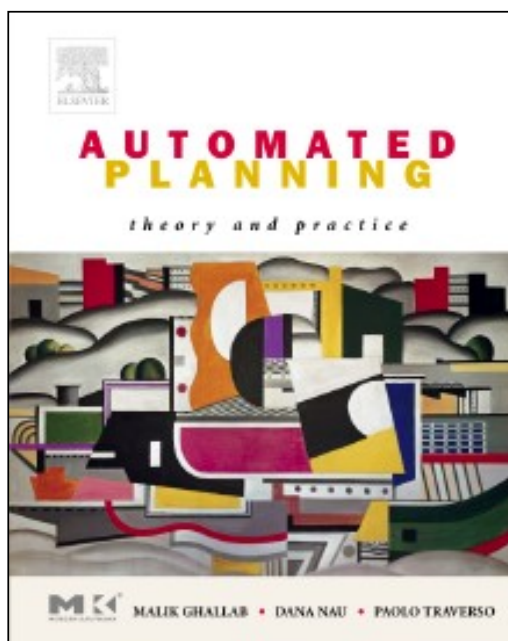
The planning algorithm will modify the control rule for next states by applying progress and if progress is false then we know that there is no plan (going through a given state) satisfying the control rule.

Forward state-space planning guided by control rules.

- If a partial plan  $S_\pi$  violates the control rule  $\text{progress}(\phi, S_\pi)$ , then the plan is not expanded.



- What we did not cover:
  - State-variable representation
  - Problem solving by transformation to SAT/CSP
  - Hierarchical task networks
  - Planning with time and resources
  - Planning with uncertainty and dynamic worlds
- What we have learned:
  - Formalization of planning problems
  - Mainstream solving approaches



## Automated Planning: Theory and Practice

- M. Ghallab, D. Nau, P. Traverso
- <http://www.laas.fr/planning/>
- Morgan Kaufmann



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