

Hypercube problems

Lecture 4

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In this lecture we list several characteristic properties of hypercubes.

1 Preliminaries

Let $G = (V, E)$ be a graph. We denote $E(A, B) := \{ab \in E \mid a \in A, b \in B\}$.

Definition 1 G has $(0, 2)$ -property if every two vertices share exactly 0 or 2 common neighbors. We also say that G is a $(0, 2)$ -graph.

An example of a $(0, 2)$ -graph other than Q_n is K_4 .

Definition 2 An interval between vertices u and v is

$$I(u, v) = \{z \in V \mid d(u, z) + d(z, v) = d(u, v)\},$$

that is, the set of vertices each of which is on some shortest (u, v) -path (geodesic).

- Every interval in Q_n forms a subcube. Specifically, $I(u, v)$ is the vertex set of $Q_n[s] \simeq Q_{d(u,v)}$ where $s = (s_1, \dots, s_n)$ is given by

$$s_i = \begin{cases} * & \text{if } u_i \neq v_i, \\ u_i & \text{otherwise.} \end{cases}$$

Definition 3 G is spherical if for every vertices u, v and $x \in I(u, v)$ there is $y \in I(u, v)$ such that $I(u, v) = I(x, y)$.

Such vertex y is unique and is called the *antipodal* vertex to x in $I(u, v)$.

- Q_n is *spherical*. The vertex y antipodal to x in $Q_n[s]$, is given by

$$y_i = \begin{cases} x_i & \text{if } s_i \neq *, \\ \bar{x}_i & \text{if } s_i = *. \end{cases}$$

Definition 4 A set $S \subseteq V$ is convex if $I(u, v) \subseteq S$ for every $u, v \in S$.

- In Q_n , convex sets \approx intervals \approx (vertex sets of) subcubes.
- There are graphs containing non-convex intervals (e. g. $K_{2,3}$).

Definition 5 The antipodal projection (*antiprojection*) of $x \in V$ into $S \subseteq V$ is

$$AP(x, S) := \{y \in S \mid d(x, y) \geq d(x, s) \text{ for every } s \in S\},$$

that is, the set of furthest vertices from x in S .

- In Q_n , any antipodal projection into any interval consists of a single vertex.

Definition 6 The median set of vertices $u, v, w \in V$ is

$$Med(u, v, w) := I(u, v) \cap I(v, w) \cap I(u, w).$$

A graph G is median if $|Med(u, v, w)| = 1$ for every $u, v, w \in V$; that is, every triple of vertices has a unique median.

Examples of median graphs: trees, grids, hypercubes. We will see in the next lecture that in some sense, they are all graphs “between” trees and hypercubes.

2 Characterizations of hypercubes

Let $G = (V, E)$ be a connected graph. The following statements are equivalent:

2.1 G is a hypercube (of dimension n).

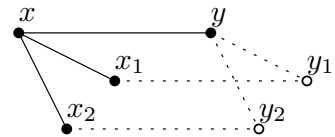
Usually, the fact that following conditions hold for hypercubes is obvious, and we just need to show the other direction (sufficiency).

2.2 G is a $(0, 2)$ -graph with $|V| = 2^n$ where n is the minimal degree. [9]

Proof Every hypercube is trivially a $(0, 2)$ -graph. The proof of the other direction is divided into three parts.

1. G is regular.

We show that every neighbor y of a vertex x has $\deg(y) \geq \deg(x)$. For every neighbor x_i of x other than y , by $(0, 2)$ -property, x_i and y have another neighbor y_i other than x . Moreover all y_i 's are distinct. From symmetry we get $\deg(x) = \deg(y)$. Since G is connected, it is also regular.



2. $|V| \leq 2^n$.

We fix some vertex x and define a level decomposition $V = L(0) \cup L(1) \cup \dots$ from x by $L(i) := \{y \in V \mid d(x, y) = i\}$. It suffices to prove the following claim.

Claim 7 For every $i \geq 0$ and $u \in L(i)$ it holds $|L(i)| \leq \binom{n}{i}$ and $|N(u) \cap L(i-1)| \geq i$.

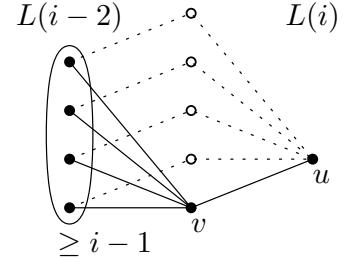
Proof By induction on i . It clearly holds for $i = 0$ and $i = 1$. Now $i \geq 2$.

The second part: every $u \in L(i)$ has a neighbor $v \in L(i-1)$. By induction for v and by $(0, 2)$ -property, we have at least $i-1$ other neighbors of u in $L(i-1)$.

The first part: from the second part we have that $i|L(i)| \leq |E(L(i), L(i-1))|$. Furthermore, by n -regularity and the second part for $L(i-1)$, we have $|E(L(i), L(i-1))| \leq (n-i+1)|L(i-1)|$. Hence,

$$|L(i)| \leq |L(i-1)| \frac{n-i+1}{i} \leq \binom{n}{i-1} \frac{n-i+1}{i} = \binom{n}{i}.$$

■



3. $|V| = 2^n$ implies Q_n .

First, $|V| = 2^n$ implies equalities in Claim 7 for every i . Thus, there are no edges within levels (G is bipartite) and the following property holds.

Property 8 In every level decomposition, every 4-cycle intersects exactly three levels.

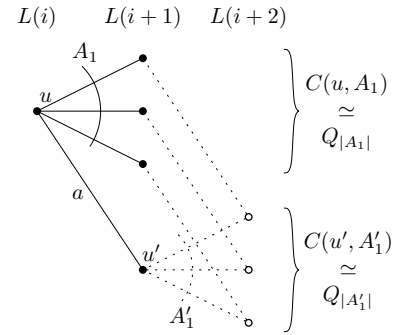
Definition 9 For $u \in L(i)$ and $A \subseteq E(u, L(i+1))$ let $C(u, A)$ be the smallest $(0, 2)$ -subgraph of G containing A .

Claim 10 $C(u, A) \simeq Q_{|A|}$ for every u and A .

Proof By induction on $|A|$. It clearly holds for $|A| \in \{0, 1, 2\}$.

Otherwise, let $A = A_1 \cup \{uu'\}$. By $(0, 2)$ -property we have a set of corresponding edges A'_1 from u' (see figure on the right). By induction, we have disjoint subcubes $C(u, A_1) \simeq Q_{|A_1|-1}$ and $C(u', A'_1) \simeq Q_{|A_1|-1}$.

Moreover, by inductively using the $(0, 2)$ -property we find a perfect matching between corresponding vertices of the subcubes, which gives us $Q_{|A|}$. ■



■

2.2' G is a $(0, 2)$ -graph with Property 8. [3]

Proof Property 8 implies Claim 7 with equalities for every i and we repeat the rest of the proof for (2.2) to imply Q_n . ■

2.3 G is n -regular, $|V| = 2^n$, and $E(N(u), N(v))$ is a perfect matching for every edge uv . [14]

Proof The last condition implies $(0, 2)$ -property. Apply condition (2.2). ■

2.4 G is bipartite, the number of shortest uv -paths is $d(u, v)!$ for every $u, v \in V$. [5]

Proof The condition implies $(0, 2)$ -property. Furthermore, bipartiteness implies Claim 7 with equalities for every i , hence $|V| = 2^n$. We finish the proof by applying (2.2). ■

2.5 $|AP(x, I(u, v))| = 1$ for every $x, u, v \in V$. [3]

That is, antiprojections into intervals are unique.

Proof

1. G is bipartite.

By contradiction, consider a shortest odd cycle C , select x arbitrarily on C and uv as the “opposite” edge to x on C . Then $AP(x, I(u, v)) = \{u, v\}$.

2. G is a $(0, 2)$ -graph.

For $ux, xv \in E$ there is exactly one common neighbor of u and v other than x , otherwise $|AP(x, I(u, v))| \neq 1$.

3. Property 8 holds.

Suppose there is a 4-cycle (u, a, v, b) in a level decomposition from x with $u, v \in L(i)$, $a, b \in L(i - 1)$ for some i . Since G is bipartite and $(0, 2)$ -graph, we have $I(u, v) = \{u, v, a, b\}$. Then we obtain a contradiction $|AP(x, I(u, v))| = \{u, v\}$.

We finish the proof by applying condition (2.2'). ■

2.5' G is $K_{2,3}$ -free and $|AP(x, C)| = 1$ for every $x \in V$ and a convex set $C \subseteq V$. [3]

Proof

1. G is bipartite. As in the previous condition.

2. G is a $(0, 2)$ -graph.

If $ux, xv \in E$, there exists another common neighbour w of u and v , otherwise $|AP(x, \{u, x, v\})| \neq 1$. A third common neighbour would form $K_{2,3}$ ($\{u, v\}$ as one partite, their common neighbours as the other one).

3. Property 8 holds. Since 4-cycles in a $(0, 2)$ -graph are convex, we can proceed as in the previous condition.

We finish the proof by applying condition (2.2'). ■

2.6 G is bipartite and every interval induces a $(0, 2)$ -graph. [11]

Proof

1. G is $(0, 2)$ -graph. If $ux, xv \in E$, then $I(u, v)$ is a 4-cycle (the only $(0, 2)$ -supergraph).
2. *Property 8 holds.* Otherwise, an interval from x (the starting vertex of a level decomposition) could contain only three vertices from some 4-cycle, which would not induce a $(0, 2)$ -graph.

We finish the proof by applying condition (2.2'). ■

2.7 $|V| = 2^n$ and $Q_0, Q_1, \dots, Q_{n-1}, G$ are the all nonisomorphic convex subgraphs of G . [15]

Proof P_3, C_3 , and $K_{2,3}$ are not convex subgraphs of G , so it is a $(0, 2)$ -graph. We finish the proof by applying condition (2.2). ■

2.8 G is bipartite and spherical. [2, 16]

Proof omitted.

2.8' G is triangle-free and spherical. [8]

Proof omitted as well.

2.9 G is bipartite and *interval-regular*. [12]

Definition 11 G is interval-regular if $|N(u) \cap I(u, v)| = d(u, v)$ for every $u, v \in V(G)$.

Proof

1. G is $(0, 2)$ -graph. If $ux, xv \in E$, then $d(u, v) = 2$ (since G is bipartite). By $|N(u) \cap I(u, v)| = 2$, there is exactly one common neighbor of u and v other than x .
2. G is n -regular (for some n). As in condition (2.2) step 1.
3. $V = 2^n$.

Take a level decomposition from x . Again, we show Claim 7 with equalities using the assumption that every $u \in L(i)$ has exactly i neighbors in $L(i - 1)$ by interval regularity for $I(u, x)$. Thus, every $u \in L(i - 1)$ is left with $n - (i - 1)$ neighbors in $L(i)$ and every $v \in L(i)$ has i neighbors in $L(i - 1)$. By induction, we have

$$|L(i)| = |L(i - 1)| \frac{n - (i - 1)}{i} = \binom{n}{i - 1} \frac{n - (i - 1)}{i} = \binom{n}{i}.$$

Consequently, $\sum_i |L(i)| = 2^n$.

We finish the proof by applying condition (2.2) (only step 3 is needed). ■

The following characterizations are listed for completeness, proofs are omitted.

2.10 G is bipartite, antipodal, and $(0, 2)$ -graph. [13]

Definition 12 G is antipodal if for every vertex u there is a vertex v such that $I(u, v) = V$.

Note that the vertex v is unique since we cannot have at the same time a shortest (u, v) -path through v' and a shortest (u, v') -path through v .

2.11 G is distance monotone, interval monotone, and $(\delta(G) \geq 3$ or $G \in \{Q_0, Q_1, Q_2\})$. [4]

Definition 13

- An interval I is closed if for every $w \in V \setminus I$ there is $w' \in I$ with $d(w, w') > \text{diam}(I)$.
- G is distance monotone if every interval in G is closed.
- G is interval monotone if every interval in G is convex.

2.12 G is bipartite, interval distance monotone, and $(\delta(G) \geq 3$ or $G \in \{Q_0, Q_1, Q_2\})$. [1]

Definition 14 G is interval distance monotone if every subgraph induced by interval in G is distance monotone.

Interval distance monotonicity is trivially implied by distance monotonicity, but not vice-versa (consider e.g. C_{2k+1}). The proof of the following characterizations will follow from the next lecture.

2.13 G is n -regular and median. [11]

2.14 G is median and $|W_{uv}| = |V|/2$ for every $uv \in E$.

Definition 15 For an edge uv we denote

$$W_{uv} := \{x \in V \mid d(x, u) < d(x, v)\}.$$

2.15 G is median and $W_{uv} = U_{uv}$ for every $uv \in E$.

Definition 16 For an edge uv we denote

$$U_{uv} := \{x \in W_{uv} \mid x \text{ has a neighbour in } W_{vu}\}.$$

2.16 G is median and for every $ab, uv \in E$, if W_{ab} and W_{uv} are disjoint, then $W_{ab} \cup W_{uv} = V$. [7]

Problems

1. Determine whether *bipartite* can be weakened to *triangle-free* (or another weaker condition) in the above characterizations.
2. Find “completely” new characterizations of hypercubes.

Notes

Characterizations (2.14) and (2.15) are proved in [6].

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