

Practical sessions for Introduction to Complexity and Computability - 4

November 19th, 2019

Exercises

- previous HWs:
 - Prove or disprove that the class of decidable (resp. partially decidable) languages is closed under symmetric difference. A symmetric difference of sets A and B is the set $A\Delta B = (A - B) \cup (B - A)$.
 - Let $D \subseteq \{0, 1\}^*$ be finite. Show that the language $\{x \mid D \subseteq L_x\}$ is partially decidable.
 - Show that the language $L_{diff} = \{\langle a, b \rangle \mid M_a(x) \downarrow, M_b(x) \downarrow \text{ and } M_a(x) \neq M_b(x) \text{ for some input } x\}$ is partially decidable but not decidable.
 - Show that $K \leq_1 Tot$ and $\overline{K} \leq_1 Tot$ where $K = \{x \mid M_x(x) \downarrow\}$ and $Tot = \{x \mid L_x = \{0, 1\}^*\}$.
 - Let $D \subseteq \{0, 1\}^*$ be finite. Show that $K \leq_1 \{x \mid D \subseteq L_x\}$ and $\overline{K} \leq_1 \{x \mid L_x \subseteq D\}$ where K is defined above.
- The “s-m-n” theorem. Find injective computable total functions $f: \mathbb{N}^2 \rightarrow \mathbb{N}$, $g: \mathbb{N} \rightarrow \mathbb{N}$ s.t.
 - $\varphi_{f(x,y)}(z) \simeq \varphi_x(z) + \varphi_y(z)$ for every $x, y, z \in \mathbb{N}$,
 - $\text{dom } \varphi_{g(x)} = \{x\}$ for every $x \in \mathbb{N}$,
 - $\text{dom } \varphi_{g(x)} = \{0, 1, \dots, x\}$ for every $x \in \mathbb{N}$.
- Rice’s theorem. Show that the following problems are undecidable.
 - Does the given program always accept (i.e. for every input)?
 - Does the given program return 0 for some input?
 - Do the given two programs compute the same arithmetic function?
- Let A be a finite problem (i.e. the language of its positive instances is finite). Determine for which functions f the class $\text{TIME}(f(n))$ contains A . Is A in the class \mathbf{P} ?
- For a language A and $k \in \mathbb{N}$ let $A_k = \{x \in A \mid |x| = k\}$ be the restriction of A to inputs of size k . Prove or disprove that for any language A , if $A_k \in \mathbf{P}$ for every k , then $A \in \mathbf{P}$.
- Consider a Turing machine M running in the following time $f(n)$ where $k \in \mathbb{N}$. In which cases the language $L(M)$ belongs to \mathbf{P} ? In which cases it belongs to \mathbf{EXP} ? Explain why.
 - $f(n) = n^k$
 - $f(n) = k^n$
 - $f(n) = n^n$
 - $f(n) = k^{k \log n}$
 - $f(n) = n^{\log \log n}$
 - $f(n) = k^{n^k}$
 - $f(n) = k^{n^n}$
 - $f(n) = n^{n^k}$
 - $f(n) = n^{n^n}$

Homework

Problems 4, 5 (both for 1 point), and 6 (all subproblems, up to 2 points).